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U.S. DEPARTMENT OF COMMERCE PATENT AND TRADEMARK OFFICE

ATTORNEY'S DOCKET NUMBER

**TRANSMITTAL LETTER TO THE UNITED STATES
DESIGNATED/ELECTED OFFICE (DO/EO/US)
CONCERNING A FILING UNDER 35 U.S.C. 371**

P-9902 MG

U.S. APPLICATION NO. (If known, see 37 CFR 1.5)

09/445386

INTERNATIONAL APPLICATION NO.

PCT/RU98/00264

INTERNATIONAL FILING DATE

July 29, 1999

PRIORITY DATE CLAIMED

November 25, 1998

TITLE OF INVENTION

Payment Method and Apparatus Therefor

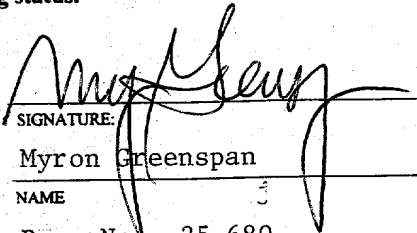
APPLICANT(S) FOR DO/EO/US Zolotorev O.A.; Kuznetsov I.V.; Moshonkin A.G.;
Smirnov A.L.; and Khamitov I.M.

Applicant herewith submits to the United States Designated/Elected Office (DO/EO/US) the following items and other information:

1. ☒ This is a **FIRST** submission of items concerning a filing under 35 U.S.C. 371.
2. ☐ This is a **SECOND** or **SUBSEQUENT** submission of items concerning a filing under 35 U.S.C. 371.
3. ☐ This express request to begin national examination procedures (35 U.S.C. 371(f)) at any time rather than delay examination until the expiration of the applicable time limit set in 35 U.S.C. 371(b) and PCT Articles 22 and 39(1).
4. ☐ A proper Demand for International Preliminary Examination was made by the 19th month from the earliest claimed priority date.
5. ☐ A copy of the International Application as filed (35 U.S.C. 371(c)(2))
 - a. ☐ is transmitted herewith (required only if not transmitted by the International Bureau).
 - b. ☐ has been transmitted by the International Bureau.
 - c. ☐ is not required, as the application was filed in the United States Receiving Office (RO/US).
6. ☒ A translation of the International Application into English (35 U.S.C. 371(c)(2)).
7. ☐ Amendments to the claims of the International Application under PCT Article 19 (35 U.S.C. 371(c)(3))
 - a. ☐ are transmitted herewith (required only if not transmitted by the International Bureau).
 - b. ☐ have been transmitted by the International Bureau.
 - c. ☐ have not been made; however, the time limit for making such amendmnts has NOT expired.
 - d. ☒ have not been made and will not be made.
8. ☐ A translation of the amendments to the claims under PCT Article 19 (35 U.S.C. 371(c)(3)).
9. ☒ An oath or declaration of the inventor(s) (35 U.S.C. 371(c)(4)).
10. ☐ A translation of the annexes to the International Preliminary Examination Report under PCT Article 36 (35 U.S.C. 371(c)(5)).

Items 11. to 16. below concern document(s) or information included:

11. ☐ An Information Disclosure Statement under 37 CFR 1.97 and 1.98.
12. ☐ An assignment document for recording. A separate cover sheet in compliance with 37 CFR 3.28 and 3.31 is included.
13. ☐ A FIRST preliminary amendment.
☐ A SECOND or SUBSEQUENT preliminary amendment.
14. ☐ A substitute specification.
15. ☐ A change of power of attorney and/or address letter.
16. ☒ Other items or information:
A Small Entity Declaration (in five signed copies)

U.S. APPLICATION NO. (if known, see 37 CFR 1.5) 09/7445386		INTERNATIONAL APPLICATION NO. PCT/RU98/0264		ATTORNEY'S DOCKET NUMBER P-9902 MG	
17. <input checked="" type="checkbox"/> The following fees are submitted: BASIC NATIONAL FEE (37 CFR 1.492 (a) (1) - (5)): \$840.00 Search Report has been prepared by the EPO or JPO \$800.00 International preliminary examination fee paid to USPTO (37 CFR 1.482) \$680.00 No international preliminary examination fee paid to USPTO (37 CFR 1.482) but international search fee paid to USPTO (37 CFR 1.445(a)(2)) \$750.00 Neither international preliminary examination fee (37 CFR 1.482) nor international search fee (37 CFR 1.445(a)(2)) paid to USPTO \$1010.00 International preliminary examination fee paid to USPTO (37 CFR 1.482) and all claims satisfied provisions of PCT Article 33(2)-(4) \$94.00				CALCULATIONS PTO USE ONLY	
ENTER APPROPRIATE BASIC FEE AMOUNT =				\$ 840.00	
Surcharge of \$130.00 for furnishing the oath or declaration later than <input type="checkbox"/> 20 <input type="checkbox"/> 30 months from the earliest claimed priority date (37 CFR 1.492(e)).				\$ -----	
CLAIMS	NUMBER FILED	NUMBER EXTRA	RATE		
Total claims	109 - 20 =	89	X \$22.00 XXXX 18.	\$ 1,602.00	
Independent claims	5 - 3 =	2	X \$78.00	\$ 156.00	
MULTIPLE DEPENDENT CLAIM(S) (if applicable)			+ \$250.00	\$ -----	
TOTAL OF ABOVE CALCULATIONS =				\$ 2,598.00	
Reduction of 1/2 for filing by small entity, if applicable. Verified Small Entity Statement must also be filed (Note 37 CFR 1.9, 1.27, 1.28).				\$ 1,299.00	
SUBTOTAL =				\$ 1,299.00	
Processing fee of \$130.00 for furnishing the English translation later than <input type="checkbox"/> 20 <input type="checkbox"/> 30 months from the earliest claimed priority date (37 CFR 1.492(f)).				\$ -----	
TOTAL NATIONAL FEE =				\$ 1,299.99	
Fee for recording the enclosed assignment (37 CFR 1.21(h)). The assignment must be accompanied by an appropriate cover sheet (37 CFR 3.28, 3.31). \$40.00 per property				\$ -----	
TOTAL FEES ENCLOSED =				\$ 1,299.00	
				Amount to be: refunded	\$
				charged	\$
a. <input checked="" type="checkbox"/> A check in the amount of \$ <u>1,299.00</u> to cover the above fees is enclosed. b. <input type="checkbox"/> Please charge my Deposit Account No. _____ in the amount of \$ _____ to cover the above fees. A duplicate copy of this sheet is enclosed. c. <input checked="" type="checkbox"/> The Commissioner is hereby authorized to charge any additional fees which may be required, or credit any overpayment to Deposit Account No. <u>10-0100</u> . A duplicate copy of this sheet is enclosed.					
NOTE: Where an appropriate time limit under 37 CFR 1.494 or 1.495 has not been met, a petition to revive (37 CFR 1.137(a) or (b)) must be filed and granted to restore the application to pending status.					
SEND ALL CORRESPONDENCE TO: Lackenbach Siegel Marzullo Aronson & Greenspan, P.C. One Chase Road, Scarsdale, NY 10583 U.S.A. Phone: 914 723 4300; Fax: 914 723 4301 December 6, 1999					
				SIGNATURE:  Myron Greenspan NAME: _____ Reg. No: <u>25,680</u> REGISTRATION NUMBER	

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RE: New National Phase PCT Application of Oleg Anatolievich
ZOLOTOREV et al. in the U.S.

For: Payment Method and Apparatus Therefor

Comprising: Complete PCT National Phase Patent Application (in
English; specification, claims, Abstract and two sheets
of drawings); Declaration and Power of Attorney;
Small Entity Declaration (five signed copies); PCT
Request Form (PCT/RO/101); Transmittal Letter to
the U.S. Designated/Elected Office; check in the
amount of one thousand two hundred ninety nine
dollars (\$1,299.00)

PCT International Application No.: PCT/RU98/00264

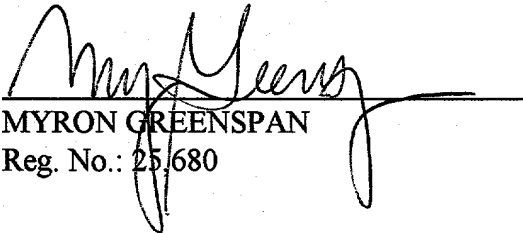
Attorney Docket No.: P-9902 MG

EXPRESS MAIL CERTIFICATION UNDER 37 CFR 1.10

"Express Mail" mailing label number EL 133 327 140 US

Date of Deposit: December 6, 1999

I hereby certify that this paper or fee is being deposited with the United States Postal Service
"Express Mail Post Office to Addressee" service under 37 CFR 1.10 on the date indicated above and
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20231.


MYRON GREENSPAN
Reg. No.: 25,680

Dated: December 6, 1999

MG/as

METHOD FOR EFFECTING PAYMENTS AND APPARATUS THEREFOR

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Field of the invention

The invention relates to electronic payment systems.

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Background of the invention

Electronic payment systems are intended to provide an adequate payment means for effecting transactions via open communication nets. Besides the degree of security and reliability, the cost of servicing, the rapidity of performing main operations, etc., an important characteristic of a payment system is the protection of user's privacy.

A user's privacy implies that nobody, not even the payment system operator, can control the user's purchases. One of the ways of protecting privacy in electronic payment systems consists in that purchases are made with the help of digital data, which confirm the solvency, but do not lead to the identification of the payer. Such data are sometimes called electronic cash. However, electronic cash, as well as any digital data, can easily be copied, so that one must take care to prevent multiple spending of electronic cash.

In certain payment systems, multiple spending is prevented by payer devices (S. Brands, Untraceable Off-Line Cash in Wallets with Observers, Advances in Cryptology CRYPTO '93, Springer-Verlag, pp. 302-318). For reliable prevention of multiple spending, such payer devices must be tamper-resistant, i.e., they must prevent unauthorized access to the data contained in the payer device. The deficiency of systems using this approach is that they are extremely unstable. The matter is that penetration into one payer device can lead to disastrous effects for the entire system, because the data contained in the payer device allow one to spend arbitrary amounts of unpaid electronic cash. Known tamper resistance technologies are not sufficiently dependable to thwart such a risk.

Electronic payment systems which do not rely on tamper resistance of payer devices must ensure, in particular, that one cannot forge payment certificates, i.e., digital data confirming the payer's solvency. The forgery is prevented by cryptographic methods, namely, by the payment system operator's digital signature. Numerous examples of digital signature are described in the books: B. Schneier, Applied Cryptography: Protocols, Algorithms, and Source Code in C, John Wiley&Sons, New York, 2nd edition, 1996 and A. J. Menezes, P. C. Van Oorshot, S. A. Vanstone, Handbook of Applied Cryptography, CRC Press, 1997.

Electronic payment systems in which the impossibility of forgery of payment certificates is ensured by the payment system operator's digital signature can be classified into offline and online ones. In offline payment systems, the moment of receiving the money by the payee is the moment of successful verification by the payee of the payment certificates provided by the payer as a payment. An advantage of such systems is that the transfer of money from payers to payees can be performed without any third party. Here, to prevent the multiple spending of payment certificates, the payer's identifier is included in payment certificates in a concealed form, and the identifier of the payer having committed multiple spending can be disclosed. The deficiency of this method is that it does not prevent the multiple spending of payment certificates, and only allows one to detect such a spending and to make a certain payer responsible for it. Thus, if the cheater is out of reach, then the payment system operator will incur losses. In addition, an honest payer's reputation can suffer if a cheater has

gained information on his payment certificates and used such a certificate. Several offline payment systems are known. For example, a system of such type is described in the patent: T. Okamoto, K. Ohta, Electronic cash system, U.S. Patent 5,224,162, 8 Jun 1992.

In an online payment system, the payee turns to the payment system operator for verification of each payment. In this case, to prevent multiple spending the payment system operator stores information on the payment certificates used earlier, and if a payment is performed with the help of a certain payment certificate, the operator checks whether the certificate has already been used before.

Known in the prior art is a method for effecting payments (Untraceable electronic cash, U.S. Patent 5,768,385, 16 Jun 1998) in which the payer receives in the bank digital signatures of payment certificates, called electronic coins, which he can use both for exchange for new electronic coins, and for payment. Here, the bank does not know which one of the two modes is used by the payer, which fact promotes untraceability of payments. Further, the multiple spending of electronic coins is prevented by the payee's online verification of the received electronic coins in the bank. However, the known method does not provide complete untraceability of a participant of the system if he is mostly a payer and not a payee, because the electronic coins given to such a participant and produced by the shop for exchange are an evidence, generally speaking, that the participant made a payment to the shop.

Known in the prior art is a method for effecting payments (D. Chaum, Security Without Identification: Transaction Systems to Make Big Brother Obsolete, Communications of the ACM, vol. 28 no. 10, October 1985 p. 1035-1038) which is the closest analog of the present invention and is chosen by the applicant as the prototype. In the known method, a client pays with payment certificates, called electronic coins, whose signatures he receives in the bank. Here, the collection of possible nominal values is fixed in advance, and for each possible nominal value of an electronic coin the bank creates a secret and a public money key. To obtain an electronic coin the payer chooses the number of the coin with the help of a random number generator, obtains the blind digital signature on the chosen number in the bank willing to credit the payer with the corresponding amount of money, and takes said digital signature as the signature of the payment certificate, which can also be called the payment certificate signature. During the payment, the payer transfers to the payee a collection of electronic coins, after which the payee verifies their validity and sends the received coins to the bank for depositing to his account. The bank verifies the validity of the electronic coins and credits the payee's account with the corresponding amount if the coins have not been already used. To control the coins already used, the bank stores the list of the numbers of used coins, the expiration dates contained in the numbers of the coins allowing the bank to delete old numbers from the list. The deficiencies of the known method are in that the bank's reputation is not defended against dishonest clients, and a client's money is not defended against a dishonest bank, because a dishonest client having received the bank's refusal to acknowledge an already used certificate for the second time can accuse the bank of cheating. In turn, a dishonest bank having received a certificate for verification may claim that the certificate has already been used before. In addition, the bank has to store information on each of the used certificates in databases with sufficiently fast access, which leads to a rapid growth of the bank's databases and to the necessity of employing expiration dates for certificates. Furthermore, in the known method the payment amount is an integral combination of nominal val-

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ues of coins, which fact either limits the range of payments, or leads to the growth of the number of the coins used in payments, which also leads to the growth of the bank's databases and slows down the payments.

Known in the prior art is an apparatus for effecting payments (T. Okamoto, K. Ohta, 5 Method and apparatus for implementing electronic cash, U.S. Patent 4,977,595, 11 Dec 1990), chosen by the applicant as the prototype.

The known apparatus for effecting payments consists of a payer device, a shop, and a bank, connected via telecommunication nets, the payer device having a means for replenishing the payer device by obtaining the blind money signature of the bank, and the bank 10 having a means for producing the money signature. In addition, the shop contains a means for offline verification of payment certificates, and the bank contains means for exposure of a cheater if he multiple-spends the bank's obligation.

The deficiency of the known apparatus for effecting payments is that it does not prevent the multiple spending of payment certificates, and only allows one to detect such a spending 15 and to make a certain payer responsible for it. Another deficiency of the known apparatus for effecting payments is that it works slowly, which is caused by the large size of the data transmitted via communication nets.

Summary of the invention

20 The main problem solved by the variants of the claimed invention is to create methods for effecting payments which would ensure an effective and reliable mechanism of paying via open communication nets, protection of each participant of the payment system from the other participants' cheating, privacy protection for ordinary participants of payments, and a wide range of payments.

25 The technical result common to all suggested variants of the claimed method for effecting payments is that when effecting payments via open telecommunication nets the financial interests of each participant are protected against the other participants' cheating, and payers and payees are able to protect their privacy. Furthermore, in some of the claimed variants, payments ranging from micro-payments to business-to-business payments are available, the 30 time taken to effect a payment depends only on the speed of action of the net connections, and not on the payment amount, and the number of the clients that can be served by the payment system operator grows proportionally to the operator's resources. The method for effecting payments is realized by an apparatus for effecting payments implemented by programming means.

35 An essential distinction between the claimed invention and the prior art is that not only the privacy of the participants of a payment operation is protected, but also the payer's financial interests are ensured, since the payment is effected on the basis of a payer order signed with the secret key connected with a payment certificate. Furthermore, in some of the claimed variants, it is allowed to spend payment certificates gradually and to replenish them.

40 The payer's privacy is protected by the procedure of making a blind signature on a payment certificate, and the payee's privacy can be protected by that when opening an account the payee does not have to give information, identifying him. In those variants where the payment system operator controls the conditions of payment and, in particular, information about the payment objective, the payer's and the payee's privacy is protected by that such a

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control is performed without access to the confidential part of the conditions of payment.

The claimed method for effecting payments is intended exclusively for hardware or computer realization, because the processing of the data used in effecting of payments, and, in particular, making and verification of digital signatures practically admits only hardware or computer realization.

The description of the method for effecting payments and the apparatus therefor as presented below is intended to describe the invention, and should not be taken to limit the scope of the claimed invention, which is described more fully elsewhere in the present specification.

10 Information specifying the terminology used in the present application is given below.

When effecting payments involving digital signatures, as in any system using digital signatures, one deals with data which are stored on suitable material media and admit digital representation (B. Schneier, *Applied Cryptography: Protocols, Algorithms, and Source Code in C*, John Wiley&Sons, New York, 2nd edition, 1996 and A. J. Menezes, P. C. Van Oorschot, S. A. Vanstone, *Handbook of Applied Cryptography*, CRC Press, 1997).

In some of digital signature schemes, it is easy to make a digital signature on random data without access to the secret keys. In order to exclude the possibility of avoid obtaining a digital signature without assistance of the holder of the secret key, one singles out among all data a set of valid data, and includes verification of their validity in verification of the validity of the signature on some data. The notion of digital signature system implies fixation of the validity criterion. In particular, coincidence of a part of the data with a sequence of bits can be the validity criterion. In another example of the validity criterion, a data is considered to be valid if it is a pair (X, Y) , where $F(X) = Y$ and F is a one-way function (B. Schneier, Applied Cryptography: Protocols, Algorithms, and Source Code in C, John Wiley&Sons, New York, 2nd edition, 1996, p. 29-30 and A. J. Menezes, P. C. Van Oorshot, S. A. Vanstone, Handbook of Applied Cryptography, CRC Press, 1997, p. 8-9), i.e., in the present case, a function which is not computationally invertible for anybody except, possibly, the signer.

Below, the keys for signing are called simply the keys, and the function of the keys used
30 for other purposes is specified.

By the payment system operator one means an entity ensuring money transactions between the participants of payment operations. In particular, the payment system operator can keep the accounts of the participants of payment operations and issue value documents. The payment system operator can consist of one bank, or it can comprise several organizations, including banks, which are interconnected by various agreements. In particular, the payment system operator's secret keys can be a secret of one of the organizations belonging to the payment system operator, and the obligations of the payment system operator to a third party can also be obligations of only one of the organizations belonging to the payment system operator. In the case where the payment system operator comprises several different banks or other organizations, there must exist a secure system of settling mutual obligations between the organizations belonging to the payment system operator. To simplify the text of the application, instead of the term "payment system operator" the term "operator" is used in what follows.

By a payment server one means an apparatus with the help of which the operator serves

effecting of payments. Such a payment server can consist of one or several computers or other devices.

By a payer device one means a device with the help of which the payer effects payments, and by a payee device one means a device with the help of which the payee effects payments. By a payment certificate one means digital data presenting the operator's obligation. A payment certificate includes base, signature, and level.

By the base of the payment certificate, also called the payment certificate base, one means data such that the signature on them, verified with the public money key, serves as a confirmation of the operator's obligations associated with this certificate; said signature is called the signature of the payment certificate or the payment certificate signature. The payment certificate signature is made with the operator's secret money key corresponding to the public money key used for verification of this signature. In what follows, the term "secret money key" is used instead of the term "the operator's secret money key".

The notion of level of a payment certificate is based on the notion of level of secret and public money keys. By the level of a public money key one means a certain numerical characteristic, associated with this public money key and defining a certain monetary value. The level can be represented, in particular, by one or several numbers. For example, the level can be represented by a nonnegative integral number L expressing in cents the monetary value defined by this level. In another example, the level can be represented by a collection of nonnegative integral numbers $L = (L_1, \dots, L_k)$, and the monetary value expressed in cents can be defined by this level by the formula $L_1 \cdot N_1 + \dots + L_k \cdot N_k$, where N_1, \dots, N_k are nonnegative integral numbers fixed in advance. In this case, the numbers L_j are called orders, and the excess of one level above another one is defined order by order. By the level of a secret money key one means the level of the corresponding public money key, and by the level of a payment certificate one means the level of the payment certificate signature, i.e., the level of that public money key with the help of which this signature can be verified.

A description of the method for effecting payments by the first variant is given below.

Effecting a payment comprises performing the payer device replenishment operation, the operation of opening the payee's account, and the payment operation per se.

One replenishes the payer device by obtaining the payment certificate signature with the help of the payment server. Here, the payment certificate signature is obtained as the blind money signature of the operator. This leads to unlinkability of the payment certificate with the replenishment source and, thus, to ensuring the payer's privacy.

One replenishes the payer device with the help of the operation of primarily filling a payment certificate, in the course of which in the payer device one creates the base of the payment certificate and obtains the payment certificate signature with the help of the payment server. A separate description of the operation of creating the base of the payment certificate is given below.

The base of the payment certificate is created in the payer device. To create the base of the payment certificate one chooses the payee's arbitrary secret key DP and the corresponding public key EP . Such a choice is performed in the framework of any digital signature system. The chosen keys DP and EP are taken as the secret and, respectively, the public key of the payment certificate. The identifier of the public key EP is included in the

base *Base* of the payment certificate. By the identifier of a public key one means arbitrary data which uniquely define the public key. In particular, the public key *EP* itself can be taken as such an identifier. In another example, the value of a cryptographic hash function (B. Schneier, Applied Cryptography: Protocols, Algorithms, and Source Code in C, John Wiley&Sons, New York, 2nd edition, 1996, p. 429) on the key *EP* can serve as an identifier. In one more example, the number of the public key received as it is registered by the operator or other entity can be taken as the identifier.

Including the identifier of the public key *EP* in the base *Base* of the payment certificate allows one to secure the payer when effecting a payment. Namely, when performing a payment operation, as described below, the operator can spend the values connected with a payment certificate only in accordance with a payer order signed by the payer. The signature on the payer order is made with the secret key *DP* and verified with the public key *EP*. Since the secret key *DP* is accessible only to the payer, such a signature also cannot be obtained without assistance of the payer.

In particular, the identifier of the public key *EP* is included in the base *Base* of the payment certificate as follows. As the base, one uses pairs (X, Y) , where $X = EP$, $Y = F(X)$, and F is a one-way function which, preferably, is collision free (B. Schneier, Applied Cryptography: Protocols, Algorithms, and Source Code in C, John Wiley&Sons, New York, 2nd edition, 1996, p. 30). In this case, one can use the data Y as the identifier of the base when performing operations with a payment certificate, and the relation $Y = F(X)$ is used as the criterion of validity of the base, where by the criterion of validity of the payment certificate base one means the validity criterion of the digital signature system to which the money keys belong.

A separate description of obtaining the signature of a payment certificate as the blind money signature of the operator during primarily filling the payment certificate is given below.

During the operation of primarily filling a payment certificate with base *Base* one obtains the payment certificate signature with the help of the procedure for obtaining the blind digital signature of the operator in the framework of an arbitrary digital signature system giving the possibility of obtaining blind digital signature.

Several digital signature systems giving the possibility of making blind digital signature are known in the prior art (D. Chaum, Blind Signature Systems, U.S. Patent 4,759,063, 19 Jul 1988; D. Chaum, Blind Unanticipated Signature Systems, U.S. Patent 4,759,064, 19 Jul 1988; D. Pointcheval, J. Stern, Provably Secure Blind Signature, Lectures Notes in Computer Science, 1163, 1996, Springer, p. 252-265, et.al.). When making a blind digital signature on initial data M , the user provides the signer with blinded data M' obtained by blinding the data M . The signer provides the user with data S' to be unblinded obtained by processing the blinded data M' with the signer's secret key, and the user obtains the signature S on the initial message by unblinding. The signature property related to the data obtained by unblinding can be verified not only after the unblinding, but also before the unblinding if the digital signature system used allows this.

When performing the operation of primarily filling a payment certificate, one takes as the initial data the identifier *BaseId* of the base of the payment certificate, i.e., arbitrary data identifying this base. When obtaining the payment certificate signature by making a blind

5 The signature on the identifier *BaseId* obtained by unblinding the data to be unblinded obtained from the payment server is accepted as the payment certificate signature with base *Base*. The validity of the made payment certificate signature can be verified with the public money key corresponding to the secret money key used when obtaining the data to be unblinded.

When performing the operation of primarily filling a payment certificate, one delivers the blinded identifier of the base to the payment server as a part of the money demand formed in the payer device, and replenishes the payer device from the funds of the replenishment source indicated in this money demand. Here, from the operator's point of view, the source of replenishment of the payer device is the source of crediting, because as a result of the replenishment operation the operator's money obligation appears in the payer device, and the payer becomes thus credited. Besides the replenishment source, the replenishment amount is also determined from the money demand if said amount is not determined by other circumstances, for example, by the terms of maintenance of the said replenishment source.

25 The operation of opening a payer's account with the operator can be performed in an arbitrary way. It is preferable that the account being opened admit a secure system of remote management.

Besides management with the help of orders signed with the secret key of the account, a public key account can also be managed by other methods. In particular, with such an account one can connect information identifying a person in charge of the account, i.e., a person having right to manage the account. Such information can be especially useful in the case of loss of the secret key of the account, because it will allow the holder to keep the control over his account in this case, too.

Moreover, if it is required to ensure the anonymity of the person in charge of the account,

then the information identifying said person can be stored in the payment server in a concealed form, which does not allow one to link this information with the person in charge of the account without assistance of a person knowing this connection. For example, this hidden information can be the value of a cryptographic hash function on the identification data of the person in charge of the account, concatenated with a password, or simply said value on the password. Information identifying the person in charge of an account can be connected with the account by an order signed with the secret key of the account when opening the account as well as at other moments of time.

A separate description of a particular case of the operation of opening a public key account is given below. This description will be used below for references in the descriptions of other variants of the claimed invention. In particular, the operation of opening an account as described below can be used when opening both the payee's account and the payer's account, and the payer's account opened in this way can be used as a source of replenishing the payer device.

The operation of opening a public key account, in particular, can proceed as follows. The future holder of the account being opened accepts an arbitrary secret key of his as the secret key of the account. The public key corresponding to the secret key of the account is delivered to the payment server and taken as the public key of the account being opened. The holder of the account considered the account to be opened after obtaining a message signed by the operator, which confirms opening the account connected with the public key.

Besides the public key of the account, other data chosen by the holder and by the operator can also be connected with the account being opened. For example, the operator can assign the account being opened a number, which is communicated to the holder of the account. It is preferable that the data which the holder connects with the account would be purely declarative, which ensures automatical performance and, thus, the cheapness of the operation of opening an account and its maintaining. The account holder's privacy can be protected by the fact that the account is opened anonymously.

A separate description of a payment operation by the first variant is given below. The descriptions of other variants of the claimed invention presented below contain references to this description. When performing a payment operation by the first variant, the value contained in the payment certificate is spent completely. Since in other variants of the claimed invention one performs, in addition to such payment operations, payment operations making it possible to gradually spend the payment certificate value, the payment operation used in the present variant is described, for more definiteness, as a payment operation with annulment of the payment certificate.

The essence of a payment operation with annulment of the payment certificate is that in the payment server one credits the payee's account, thus canceling the operator's obligation associated with the payment certificate. The payee's account is credited in the case where the payment certificate have not been utilized, the payment certificate signature is valid, and the signature on the payer order is valid. One judges whether the payment certificate has already been used by the information on the used payment certificates which is contained in the payment server information storage, while the payment certificate signature and the payer order signed with the secret key of the payment certificate is delivered to the payment server with the help of the payee device. A more detailed description of a payment operation with

annulment of the payment certificate is as follows.

When performing a payment operation, one forms in the payer device a payer order, in which one includes information about the payee and the identifier of the base of the payment certificate. The payer order signed with the secret key of the payment certificate includes the data delivered to the payee device and called the payment data. In particular, the information about the payee included in the payer order contains the conditions of payment and the identifier of the payee's account if this account is not determined by other circumstances.

By the conditions of payment one means arbitrary data specifying the circumstances of effecting a payment. Such data can include, in particular, the payment amount and restrictions on the time taken to perform a payment. In particular, the conditions of payment may contain, possibly, in a form hidden from the operator, data describing the payee's obligations implied by the fact of effecting the payment.

In the payee device, one forms data called the payee order, which are delivered to the payment server. Here, one includes in the payee order the payer order contained in the payment data, and, possibly, other data, including, in particular, the conditions of payment. In addition, the payer order can be signed with the secret key of the payee's account if the payer's account is a public key account.

In the payment server, one credits the payee's account on the basis of the payee order and forms the operator's response to the payee order, by means of which response the payment operation is judged. Here, the payee's account is credited if the payment server information storage does not contain information saying that the payment certificate has already been used. Moreover, in the payment server, the validity of the payment certificate signature and the validity of the signature on the payer order are verified before crediting the payee's account, the signatures being taken from the payer order contained in the payee order received from the payee device. To prevent the multiple spending of the payment certificate, one enters the information on said payment certificate into the payment server information storage. To confirm the rights of the operator to effect the payment operation being performed, the signed payer order is also entered into the payment server information storage. The operator's response to the payer order is delivered to the payee device.

During the execution of a payment operation, the payer's financial interests are protected by the operator's duty to verify the signature on the payer order, while the signed payer order stored in the information storage protects the operator against false accusations that he has misappropriated the value contained in the payment certificate.

The payer's privacy is protected, in particular, by that the payer order intended for delivery to the payment server is encrypted with the operator's encryption key, where by the encryption key of an entity one means a key for encrypting the messages sent for to this entity.

Besides the payer order, the payment data can include data intended for the payee, and, in particular, the identifier of the service or goods to be paid for.

In particular, the payee's security is ensured by the fact that the payee receipt signed with the operator's secret key is included in the operator's response to the payee order delivered to the payee. Here, by the payee receipt one means data confirming the fact of crediting the payee's account. Such a receipt, besides the amount of crediting and the identifier of the credited account, may also contain other data, in particular, the conditions of payment, the time taken to perform the crediting operation, etc. Having verified the validity of the opera-

tor's signature on the payee receipt, the payee regards the payment as being performed and delivers to the payer data confirming the fact of the performing of the payment.

Moreover, during the execution of a payment operation, in the payee device, by means of the operator's response to the payee order, one forms and delivers to the payer device data according to the performing of the payment for payer is judged. These data can contain, in particular, both the payee's consent to the fact that the payment operation has been successfully performed for the payee, and the payer receipt signed with the operator's arbitrary secret key. In the payment server, such a receipt is included in the operator's response to the payee order sent to the payee, and in the payee device the receipt is taken from the operator's response to the payee order received by the payee. Here, by the payer receipt one means data confirming the fact of spending the value contained in the payment certificate during the execution of a payment operation. Furthermore, the payer receipt may contain other data and, in particular, the conditions of payment and the time taken to perform a payment by the operator. In addition, before including the payer receipt into the operator's response to the payee order, one can encrypt said receipt with an arbitrary encryption key of the payer. Moreover, along with the data according to which the performing of the payment for the payer is judged, one can deliver to the payer device data connected with the fulfilment by the payee of his obligations implied by the fact of payment, and, in particular, the password of access to certain information.

During a payment operation, the value contained in the payment certificate can be spent either completely on the payment objectives, or partially on other objectives. In particular, a part of the payment certificate value can be returned to the payer device as change. A separate description of such an operation is given below.

When returning a part of the payment certificate value, one replenishes the payer device either by primarily filling the payment certificate or by replenishing the payment certificate already at hand. Moreover, the returned value of the payment certificate, i.e., the unspent part of the payment certificate value, can be returned in such a way that the returned value will remain unknown to the operator. Such a reimbursement of the returned value can be performed, for example, with the help of the method described in the patent: David Chaum, Returned value blind signature systems, U.S. Patent 4 949 380, 14 Aug 1990. In the best embodiment of the operation of returning a part of the payment certificate value to the payer device, the operator cannot distinguish primarily filling the payment certificate from replenishing a payment certificate already at hand, which promotes the protection of the payer's privacy.

The payee's security during the execution of a payment operation is ensured, in particular, by the fact that the payee receipt signed with the operator's arbitrary secret key is included into the operator's response to the payee order, while the performing of the payment for the payee is judged according to the validity of the signature on the payee receipt.

Furthermore, the payer can control the circumstances of performing a payment operation. For such a control, one includes the conditions of payment into the payer order when forming the payer order. In particular, the conditions of payment contained in the payer order, can include payee obligation data, i.e., data describing the payee's obligations implied by the fact of effecting the payment. To ensure that in case of need the payer could insist on the fulfilment by the payee of his obligations, the payee obligation data are signed with the payee's

arbitrary secret key, and before the payer decides to perform a payment operation, one verifies in the payer device the validity of the signature on the payee obligation data and stores the signed payee obligation data. If it is required to conceal the payment objective from the operator, then in the payer device, when forming the payment data, one processes the payee obligation data with the help of an arbitrary one-way function, and includes the data obtained by the processing in the payee order as the conditions of payment.

In addition, during the execution of a payment operation the payee can play a role of a payer. Such a payment operation can be used for transferring values from the payer device to the payer's account.

Note a particular case of performing the payer device replenishment operation, in which the replenishment is performed from the funds of an intermediate payer. In particular, for such performing of the operation of replenishment of the payer device, the data blinded in the payer device during obtaining the blind money signature of the operator are subjected to an additional blinding in an intermediate payer's payer device, and the data to be unblinded that were received from the operator are subjected to the corresponding additional unblinding. The additional blinding is intended for the protection of the privacy of the intermediate payer.

A common problem of all electronic queueing systems caring about the users' privacy is connected with potential possibility of identifying the user by net address information. Several theoretical approaches to the solution of this problem are described in the paper: D. Chaum, Security Without Identification: Transaction Systems to Make Big Brother Obsolete, Communications of the ACM, vol. 28 no. 10, October 1985 p. 1031-1034. A practically acceptable method of concealing the user's net address information from the addressee of the message can be to use remailers transferring the message to the addressee from their own address. In particular, in all variants of the claimed invention during the performing of a payment operation, the payer's address is, to a certain degree, hidden from the bank by that the payer order is delivered to the payment server via the payee device.

In some embodiments of the method for effecting payments by the first variant, payments in the range from micro-payments to business-to-business payments are possible, and the time taken to perform a payment depends only on the speed of action of the net connections and not on the payment amount, which is achieved by choosing levels of payment certificates such that the value defined by the level can take arbitrary values in the mentioned range, and by effecting a payment operation with the help of a beforehand-bounded number of payment certificates of suitable value, whose number is bounded in advance, or even a single payment certificate.

A description of the method for effecting payments by the second variant is given below.

One effects payments by the second variant in the same way as by the first variant, with the only exception, that besides replenishing the payer device with the help of primarily filling the payment certificate one also replenishes the payer device by replenishing the payment certificate. Further, in the best embodiment, when replenishing the payer device, the operator cannot distinguish primarily filling the payment certificate from replenishing a payment certificate already at hand, and the source of funds for such a replenishment is in no way connected with the source of funds for primarily filling.

If, when performing a payment operation by the first variant, one presents to the payment server the signature of a payment certificate whose level is the maximal available in the payer device, then the operator can detect the coincidence of this level with the level of the money keys used when replenishing the payer device, which can lead to linking the payment certificate used in a payment with the source of its replenishment. The probability of this is especially high in the case where the possible levels of payment certificates are such that the value defined by the level can take arbitrary values in the range from micro-payments to business-to-business payments. Replenishing the payer device by replenishing the payment certificate reduces the possibility of such linking considerably, because after replenishing the payment certificate, its level is the sum of several levels of signature with secret money keys. Thus, replenishment of the payment certificate additionally promotes the protection of the payer's privacy. In addition, replenishment of the payment certificate additionally allows one both to accumulate on a single payment certificate the returned values of used payment certificates reimbursed in payment operations, and to replenish said payment certificate from the funds of another source.

A separate description of the operation of replenishing the payer device with the help of replenishing the payment certificate is given below.

The payment certificate replenishment operation, i.e., the operation of obtaining a signature of the payment certificate, whose level exceeds the level of that payment certificate signature which is contained in the payer device at the beginning of the replenishment operation, is performed with the help of obtaining the blind money signature of the operator. Here, one takes as the initial data the payment certificate signature, which is contained in the payer device. For performing the payment certificate replenishment operation, the correspondence between the levels and the money keys must be such that the signature of level A on certain data X coinciding with the signature of level B on certain data Y corresponds to the signature of level $(A + B)$ on the data Y . An example of such correspondence is the case where the level L is represented by a collection of nonnegative integral numbers (L_1, L_2, \dots, L_k) , and for each index j from 1 to k there is a function S_j , whose calculation requires the secret key K_j . In the case where the functions S_j commute with each other, i.e., $S_j(S_i(X)) = S_i(S_j(X))$ for an arbitrary data X , arbitrary data allowing one to calculate the function S_L , which is the composition of the functions S_j , each of which occurs in the composition with multiplicity L_j , can be taken as the money key corresponding to the level L . In this case, the data $S_L(X)$ are the signature on the data X with the money key corresponding to the level L . For example, such a key can be represented by a collection of data $(L, K_1, K_2, \dots, K_k)$. An example of a specific system of such functions allowing an economical making and storage of secret keys is presented below in the description of the best embodiment of the method for effecting payments. Note that the operation of replenishing a payment certificate is a particular case of the payment certificate replenishment operation, if the identifier of the base of the payment certificate is taken as the payment certificate signature of zero level.

In the best embodiment of the operation of replenishing payment certificates by means of increasing their level with the help of the operator making the blind payment certificate signature of higher level, the operator has no possibility to determine whether the payer replenishes a payment certificate already at hand, or he performs the operation of filling a newly

A description of the method for effecting payments by the third variant is given below.

One effects payments by the third variant in the same way as by the first variant, with the only exception that one uses payment certificates which allow one to perform several payment operations. This is reflected in the fact that the operator opens a payment account connected with the given payment certificate and connects said account with the public key of the payment certificate signature. Further, one credits the payee's account during the execution of a payment operation from the funds of the payment account, which is credited from the funds of the payment certificate. Thus, it is possible to gradually spend the value contained in the payment certificate.

A separate description of the operation of opening a payment account is given below.

30 A separate description of the operation of crediting a payment account is given below.

A separate description of the payment operation used in the method for effecting payments by the third variant is given below. In distinction to the payment operation with annulment of the payment certificate as described earlier, the payment operation used in the method for effecting payments by the third variant is described as a payment operation in-

A payment operation involving a payment account is performed in the same way as the above-described payment operation with annulment of the payment certificate, with the exception of the following. First, one credits the payee's account from the funds of the payment

40 Another advantage of performing the payment operation involving a payment certificate in comparison to the same payment operation with annulment of a payment certificate, in which one uses a collection of payment certificates, is that the time taken to effect a payment depends only on the speed of action of the net connections, and not on the payment amount, because the fixed size of payer orders makes it possible to indicate any practically possible

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amount of money as the payment amount.

In addition, a payment account connected with one of the payment certificates can be used in the payer device replenishment operation as the replenishment source. Funds can thus be transferred from the payment account to the payer device.

5 A description of the method for effecting payments by the fourth variant is given below.

One effects payments by the fourth variant in the same way as by the third variant, with the only exception that as in the method for effecting payments by the second variant described above one additionally performs the payment certificate replenishment operation. Additional advantages of performing such an operation are also indicated in the above description of the method for effecting payments by the second variant.

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The description of the apparatus for effecting payments is given below. Using this apparatus for effecting payments, it is possible to realize the claimed method for effecting payments by the third and fourth variants.

The apparatus for effecting payments contains a payer device, a payee device, and a payment server connected with telecommunication nets. Each of these devices can be realized by a computing device programmed in the corresponding way.

Such a computing device can be chosen from a large number of known electronic devices, for example, personal computers. Such computing devices include, in an internal or an external way, storage devices, which are intended for storing data or codes of programs that are involved in effecting payments. In addition, such computing devices comprise auxiliary devices, for example, modems, which enable the computing devices to communicate with other similar devices. The communication environment, in the framework of which the data are exchanged, can also be any of a large number of possibilities, including telephone lines, cable, the Internet, satellite or radio transmissions, optical-fiber connections, etc. In other words, it is not assumed that the invention is limited with regard to either the types of devices that are used or the methods of communication that are employed.

The payer device, the payee device, and the payment server can be realized by the wide variety of programming means based on the corresponding algorithms. Here, the payer device contains a means for replenishing the payer device by the use of making a blind money signature of an operator, a means for creating a payment certificate base by processing a public key of the payment certificate with a one-way function, a means for storing the created payment certificate base in a storage device, and a means for forming a payer order signed with a secret key of the payment certificate. Here, the means for replenishing the payer device by the use of making a blind money signature of an operator is realized by a means for increasing the level of the payment certificate signature.

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Said means for increasing the level of the payment certificate signature has a means for forming a money demand including a blinded payment certificate signature, a means for unblinding the data to be unblinded comprised in a response to the money demand, and a means for entering the result of unblinding into said storage device.

40 The means for increasing the level of the payment certificate signature is intended for replenishing the payer device either by primarily filling a payment certificate or by replenishing a payment certificate. In the former case, with the help of the means for increasing the level of the payment certificate signature one processes the payment certificate signature of zero level, which coincides with the identifier of the payment certificate base created with

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In addition, the payee device contains a means for opening a public key account.

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Brief description of the drawings

In what follows, the present invention is clarified by the description of specific examples of its implementation and by accompanying drawings, where:

- Fig.1 shows a flow-chart of the apparatus for effecting payments;
- 5 Fig.2 shows a flow-chart of the payer device replenishment operation;
- Fig.3 shows a flow-chart of the payment operation.

Best embodiment of the invention

In the best embodiment of the method for effecting payments by each of the variants, payments are effected in the framework of a payment system, whose operator comprises a variety of banks, and the system has many payers and payees. Further, both the payers and the payees utilize devices simultaneously containing both a payer and a payee device. Furthermore, in the description of the best embodiment of the invention, such a device is called the "Electronic wallet".

15 Before serving a client, a bank belonging to the payment system performs preparatory actions. Since the preparatory actions are performed in the same way in the best embodiment of each of the variants of the claimed method, a separate description of the best way of performing of these operations is given below.

At the stage of preparatory actions, one fixes a digital signature system giving a possibility of making blind digital signature. This system is intended for making and verification of money signature, and below it is called a money signature system. One also fixes a collection of admissible levels, i.e., quantities each of which defines a certain monetary value. Here, the collection of admissible levels is chosen so that an arbitrary monetary value presenting a practical interest when replenishing the payer device would be represented by some level. For each admissible level, the bank chooses a secret money key corresponding to this level and a public money key corresponding to the secret money key in the framework of a fixed money signature system. Here, the secret money key corresponding to each of the admissible levels is chosen so that the signature of level A on some data X , X coinciding with the signature of level B on certain data Y , corresponds to the signature of level $(A + B)$ on the data Y . Information on the public money keys and the monetary values corresponding to admissible levels is made public and is entered into the storage devices of "Electronic wallets".

One also fixes a digital signature system intended for signing the messages used when effecting payments. In the framework of this system, the bank chooses secret keys and the corresponding public keys. The information on public keys is made public and is entered into the storage devices of "Electronic wallets".

In addition, one fixes the structure and the validity criterion of payment certificate bases, as well as the way of including the identifier of the public key of the payment certificate into the base of the payment certificate. To this end, one fixes a cryptographic hash function F which takes values in bit strings and, preferably, is collision free. Payment certificate bases are pairs (Y, X) , where Y is the public key of the payment certificate, which key is chosen in the framework of a fixed signature system and used as its own identifier. Further, a base with such a structure is considered to be valid if $F(Y)=X$. One takes the data X as the identifier *BaseId* of the base of the payment certificate.

The possibility of realization of the preparatory actions described above in the best embodiment of the method for effecting payments by each of the claimed variants is clarified by the following example.

Example 1

- 5 As the money signature system, one uses a system based on the digital RSA-signature (B. Schneier, Applied Cryptography: Protocols, Algorithms, and Source Code in C, John Wiley&Sons, New York, 2nd edition, 1996 and A. J. Menezes, P. C. Van Oorshot, S. A. Vanstone, Handbook of Applied Cryptography, CRC Press, 1997). In this system, a secret key is the pair (N, D) , where D is a secret exponent and N is the modulus. Here, the corresponding public key is a pair (N, E) , where E is a public exponent satisfying the condition $X^{E \cdot D} \equiv 1 \pmod{N}$ for each integral number X coprime to N . The RSA system allows several methods of making a blind digital signature. For example, one of these methods is described in the patent: D. Chaum, Blind Signature Systems, U.S. Patent 4,759,063, 19 Jul 1988, and another one is described in the patent: D. Chaum, Blind Unanticipated Signature Systems, 15 U.S. Patent 4,759,064, 19 Jul 1988.

As an admissible level, one takes an arbitrary collection of nonnegative integral numbers $L = (L_1, L_2, L_3)$, and such a level defines the monetary value expressed in cents by the formula $L_1 \cdot N_1 + L_2 \cdot N_2 + L_3 \cdot N_3$, where $N_1 = 100^2$, $N_2 = 100$, $N_3 = 1$.

- For each admissible level, the corresponding secret money key and the public money key 20 corresponding to the secret money key are chosen in the framework of the RSA-signature system as follows. As the modulus of each of the public and secret money keys, one uses an arbitrary RSA-modulus N , for which the integral numbers $E_1 = 3$, $E_2 = 17$, $E_3 = 5$ are admissible as public exponents. Below, the numbers E_1, E_2, E_3 are called the basic public exponents. One takes the key with public exponent $E = E_1^{L_1} \cdot E_2^{L_2} \cdot E_3^{L_3}$ as the public money key 25 corresponding to the level L . As the secret money key corresponding to the level L one takes the key with secret exponent $D = D_1^{L_1} \cdot D_2^{L_2} \cdot D_3^{L_3}$, where D_j is the secret exponent corresponding to the basic public exponent E_j . Means for creating keys in the RSA system are well known (B. Schneier, Applied Cryptography: Protocols, Algorithms, and Source Code in C, John Wiley&Sons, New York, 2nd edition, 1996 and A. J. Menezes, P. C. Van Oorshot, S. A. Vanstone, Handbook of Applied Cryptography, CRC Press, 1997). The modulus N and basic public exponents E_1, E_2, E_3 are published as the information about public money keys. 30 The data $N_1 = 100^2$, $N_2 = 100$, $N_3 = 1$ are published as the information on the monetary values corresponding to the admissible levels.

- The cryptographic hash function F used of the criterion of validity of payment certificate 35 bases, is chosen so that its value on the bit sequence X is obtained by concatenation of bit sequences $H(X)$ and Y , where H denotes the known hash function SHA-1 (A. J. Menezes, P. C. Van Oorshot, S. A. Vanstone, Handbook of Applied Cryptography, CRC Press, 1997, p. 348), $Y = 1111...1110$, and the number of the unit bits in Y is such that the total number of bits in $F(X)$ is equal to the number of bits in the modulus N .

- 40 As the digital signature system intended for signing the messages used when effecting payments, one fixes the RSA system with public exponent 3, i.e., the RSA system in which the public exponent is fixed and equal to 3. In the framework of this system, the bank chooses a secret key DB and the corresponding public key EB . The modulus of the public key EB is made public as the information about this key.

A description of the best embodiment of the method for effecting payments by the first variant is given below.

Having performed the preparatory actions described above, the payee opens a public key account in the bank. Here, the account is opened in the following way.

- 5 As the secret key of the account being opened, the payee takes his secret key created specially for this purpose in the framework of the signature system fixed at the stage of preparatory actions. The public key corresponding to the secret key of the account is delivered to the payment server via open telecommunication nets. In the payment server, one takes the delivered public key as the public key of the account being opened, assigns a number to the
- 10 account being opened, and creates in the storage of accounts a record containing the number of the account, the public key of the account, and other attributes of the account. The payee considered the account to be opened after receiving a message signed with the bank's secret key, which confirms opening an account connected with the public key of the account. Further, the signature on the message is stored in the storage device of the "Electronic wallet",
- 15 which will allow one in the future to make justified claims to the bank if the bank has not fulfilled its obligations connected with the opened account.

- Advantages of the described variant of opening an account are that the operation of opening an account proceeds in a remote manner in the real-time mode. In addition, opening an account in the payment server proceeds automatically, because in particular, it requires no
- 20 verification of the data identifying the account holder. Thus, this variant of opening an account is very economical for the bank and very convenient for the bank's clients. In addition, as every public key account, an account opened as described above admits a secure remote management, and a relatively small size of the number of the account allows one to reproduce it by hand, manually in particular, for indicating the designation of a postal money order.

- 25 In the best embodiment, the payer uses an arbitrary means for money operations which allows a secure remote management as a source of replenishment of his "Electronic wallet".

- When creating the payment certificate base *Base* in the course of performing the operation of primarily filling a payment certificate, as the payee's secret key *DP* and the corresponding public key *EP* one chooses keys specially created for this purpose. Here, the keys
- 30 *DP* and *EP* are chosen in the framework of the signature system fixed at the stage of preparatory actions.

- During the operation of primarily filling a payment certificate, the blinded identifier *BaseId* of the base of the payment certificate is delivered to the payment server as a part of a money demand. Here, the data necessary for remote withdrawal of funds from the utilized
- 35 replenishment source are also included in the money demand.

- Prior to performing a payment operation, the payee obligation data are signed with the payee's arbitrary secret key, and before the payer decides to perform a payment operation, in the payer device one verifies the validity of the signature on the payee obligation data and
- 40 stores the signed payee obligation data.

The information about the payee included in the payer order when performing a payment operation contains the identifier of the payee's account and the conditions of payment. The result of processing the payee obligation data with the help of a fixed one-way function is included in the conditions of payment. The payer order is encrypted with the operator's en-

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ryption key.

A session key for symmetric encryption is included in the payee order and the payee order itself is encrypted with the operator's encryption key.

The payee receipt signed with the operator's secret key is included in the operator's response to the payee order, and the operator's response to the payee order is encrypted with the session key contained in the payee order. The payee judges the performing of the payment according to the validity of the signature on the payee receipt.

In the payee device, by means of the operator's response to the payee order, one forms and delivers to the payer device the data confirming the payee's consent to the fact that the payment operation has been successfully performed.

The part of the payment certificate value that was not spent on the payment objectives is returned to the payer device as change.

The possibility of realization of the above-described best embodiment of the method for effecting payments by the first variant is clarified by the following example.

15 Example 2

The preparatory actions are performed as in Example 1 described above.

The payee, who is a seller in the present example, opens an account with public key *ES* in the bank as described above. Further, in the storage device of the seller's "Electronic wallet" one creates a record with information about the opened account, which information contains the secret key *DS* corresponding to the public key *ES*, the number *AccountId* of the opened account, and other attributes of the account.

In this example, the source of replenishment of the payer's "Electronic wallet" is the payer's account, which he opens in exactly the same way as the payee. The payer delivers a certain amount of money to the account opened by him by a postal money order; for the sake of definiteness, let it be 100 dollars.

Using the opened account as the replenishment source, the payer replenishes his "Electronic wallet" with the help of the operation of primarily filling a payment certificate.

During the operation of primarily filling a payment certificate, the base of the payment certificate is created by choosing the secret signature key *DP* and the corresponding public signature key *EP* in the framework of the RSA system with public exponent 3. Since the public exponent is fixed, the public key is represented by a single modulus. As indicated above, means for creating such keys are well known. One takes the data (EP, X) as the base (Y, X) of the payment certificate, where *X* is obtained from the key *EP* with the help of the means for calculating the function *F* described in Example 1 presented above. Such a means is realized by programming means on the basis of the above description for the function *F*.

During the operation of primarily filling a payment certificate, the blinded identifier *BaseId* of the payment certificate base is delivered to the payment server as a part of the money demand. One also includes in the money demand the required replenishment amount, for the sake of definiteness 200 dollars, and the payer's account number, and along with the money demand one delivers to the payment server the signature of said demand with the secret key of the payer's account. Here, the payer device is replenished from the funds of the account indicated in the money demand, after verification of the signature contained in the money demand with the help of the public key of the account indicated in

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To blind the identifier *BaseId* of the payment certificate base one chooses a level $M = (M_1, M_2, M_3)$ such that L_1 does not exceed M_1 , L_2 does not exceed M_2 , and L_3 does not exceed M_3 , where $L = (L_1, L_2, L_3)$ is the level of that secret money key which will be used in the payment server when replenishing the payer device. The level M to be chosen depends on the required replenishment amount expressed in cents, i.e., equal to 20 000 in the present example. When choosing the level M , one assumes that the level L of the secret money key used by the bank is defined as follows according to the rule fixed in advance in the payment system.

In any case, the fixed rule of defining the level L in accordance with the required replenishment amount guarantees, in this example, that L_1 does not exceed 2, while L_2 and L_3 do not exceed 99. Thus, one takes the data (2, 99, 99) as the level M .

In the payment server, the signature on the blinded data X' with the secret money key corresponding to the level L is taken as the data S' to be unblinded. Thus, in this example, the data X' is processed with the secret money key with modulus N and secret exponent $D = D_1^{L1} \cdot D_2^{L2} \cdot D_3^{L3}$. After that, the data S' are delivered to the payer device.

The payer willing to pay 43.50 dollars to the seller for some goods, prepares the pay-

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ment data *PaymentData*, including in them the payer order *PayerOrder* signed with the secret money key of the payment certificate *DP* and data *A* intended for the seller and consisting in this example of the name of the paid goods and the identification data of the recipient of the goods. The payer order *PayerOrder* consists of the public key of the payment certificate *EP*, the signature *S* of the payment certificate, the seller's account number *AccountId*, and the data *C* defining the conditions of payment. In the present example, the payer takes as *C* the seller's account number *AccountId* and the value of the hash function *H* on the text of the seller's obligation valid in the case of a payment, namely, an obligation to supply the corresponding goods to the person with the indicated identification data.

10 The seller willing to accept the payment forms the payee order *SellerOrder* = (*AccountId*, *PayerOrder*) encrypted with the bank's public encryption key and delivers it to the bank.

Having made sure that the list of the used payment certificates contains no record on the payment certificate with public key *EP*, having verified the validity of the signature on the payer order *PayerOrder* with the help of the public key *EP*, and having verified the validity of the signature *S* of the payment certificate, the bank enters a record including the public key *EP* and the signed payer order *PayerOrder* into the list of used payment certificates and credits the account with number *AccountId* with an amount of 43.49 dollars, under the assumption that the bank charges 1 cent for performing a payment operation. After that, the bank forms the payer receipt confirming the fact of crediting the account with number *AccountId* with an amount of 43.49 dollars, signs it, and delivers to the seller, who, having verified the validity of the bank's signature on the receipt received, regards the payment as being effected and informs the payer about successfully effecting the payment.

25 The rest value of the payment certificate, making quantity (19 915 - 4 350) cents in the present example, is returned by replenishing the payer device, and the replenishment can be performed either by primarily filling the payment certificate or by replenishing a payment certificate already at hand. To this end, one includes in the payer order a blinded signature of that payment certificate whose value will increase due to the returned value, while the corresponding money signature of the operator, delivered to the payer device through the payee device along with the other data is included in the operator's response to the payee order.

A description of the best embodiment of the method for effecting payments by the second variant is given below.

35 In the best embodiment of the method for effecting payments by the second variant, one performs all actions that are performed in the best embodiment of the method for effecting payments by the first variant.

Moreover, one performs an additional payment certificate replenishment operation, which is realized so that the operator cannot distinguish the operation of primarily filling a payment certificate from the payment certificate replenishment operation. One performs the payment certificate replenishment operation at an arbitrary moment of time before the payment operation and one uses for replenishing a payment certificate whose value is not sufficient of the payment operation.

The possibility of realization of the above-described best embodiment of the method for effecting payments by the fourth variant is clarified by the above-described Example 2,

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where the possibility of realization of the best embodiment of the method for effecting payments by the first variant is clarified, and by the following Example 3, where the possibility of realization of the best embodiment of replenishing the payer device with the help of replenishing the payment certificate is clarified.

5 Example 3

In this example, we use agreements and notation adopted in Example 1 described above.

The payment certificate replenishment operation, i.e., the operation of obtaining a payment certificate signature whose level exceeds the level of the payment certificate signature contained in the payer device at the beginning of the replenishment operation, is performed by obtaining the blind money signature of the operator. Here, as the initial data one takes the payment certificate signature contained in the payer device.

Suppose that the payer's "Electronic wallet" contains a payment certificate with signature S_I of level $(0, 2, 30)$. Suppose that the value of this payment certificate, which is equal to 2 dollars and 30 cents in the present case, is not sufficient for the payer to make a certain purchase. In order to make this purchase anyway, the payer replenishes his "Electronic wallet" by replenishing the payment certificate.

For this purpose, one delivers the blinded signature S_I of the payment certificate to the payment server as a part of the money demand. One also includes in the money demand the required replenishment amount, equal, for the sake of definiteness, to 50 dollars, and the payer's account number, and along with the money demand one delivers to the payment server the signature of said demand with the secret key of the payer's account. Here, the payer device is replenished from the funds of the account indicated in the money demand after verification of the signature contained in the money demand with the help of the public key of said account.

To blind the signature S_I of the payment certificate, as in Example 2 presented above, one chooses a level $M = (M_1, M_2, M_3)$ such that L_1 does not exceed M_1 , L_2 does not exceed M_2 , and L_3 does not exceed M_3 , where $L = (L_1, L_2, L_3)$ is the level of that secret money key which will be used in the payment server when replenishing the payer device. One chooses the level M depending on the required replenishment amount as in Example 2. In the present example, $M = (0, 5, 99)$.

As in Example 2, one defines the replenishment amount by the required amount contained in the money demand and by the funds on the payer's account. Suppose that the replenishment amount in this example turns out to be equal to 47 dollars and 13 cents, i.e., 4 713 cents. One defines the level L by the replenishment amount as in Example 2. In the present example, $L = (0, 47, 13)$.

After the level M is chosen, the initial data X , which is taken equal to the identifier $BaseId$ of the base of the payment certificate during primarily filling the payment certificate, is blinded in accordance with the relation $X' = F \cdot X \pmod{N}$, where $F = R^U \pmod{N}$, $U = U_1 \cdot U_2 \cdot U_3$, $U_1 = E_1^{M_1}$, $U_2 = E_2^{M_2}$, $U_3 = E_3^{M_3}$, and R is a randomized integral number of suitable size.

In the payment server, the signature on the blinded data X' with the secret money key corresponding to the level L is taken as the data S' to be unblinded. Thus, in this example, the data X' is processed by the secret money key with modulus N and secret exponent $D = D_1^{L_1} \cdot D_2^{L_2} \cdot D_3^{L_3}$. After that, the data S' are delivered to the payer device.

In the payer's "Electronic wallet", from the received data S' to be unblinded one makes the signature S of the payment certificate by unblinding the received data S' in accordance with the relation $S = S' \cdot T^{-1} \pmod{N}$, where $T = R^V \pmod{N}$, $V = V_1 \cdot V_2 \cdot V_3$, and $V_1 = E_1^{M1-L1}$, $V_2 = E_2^{M2-L2}$, $V_3 = E_3^{M3-L3}$. The made signature S is stored in the storage device instead of the signature S_1 stored before. Here, the level of the signature S is equal to the sum of the level of the signature S_1 , i.e., (0, 2, 30) in the present example, and the level $L = (0, 47, 13)$. Thus, the level of the signature S is equal to (0, 49, 43), and the payment certificate value has increased to 49 dollars and 43 cents.

A description of the best embodiment of the method for effecting payments by the third variant is given below.

The preparatory actions, opening the payee's account, replenishment of the payer device, and the payee's actions during the execution of a payment operation are performed as in the above-described best embodiment of the method for effecting payments by the first variant.

As in the above-described best embodiment of the method for effecting payments by the first variant, before the payer decides to perform a payment operation, in the payer device one verifies the validity of the signature on the payee obligation data and stores the signed payee obligation data.

When performing a payment operation, the identifier of the payee's account and the conditions of payment are included in the payer order as information about the payee. One includes in the conditions of payment the result of processing the payee obligation data with the help of a fixed one-way function. The payer order is encrypted with the operator's encryption key.

A session key for symmetric encryption is included in the payee order, and the payee order itself is encrypted with the operator's encryption key.

The payee receipt signed with the operator's secret key is included in the operator's response to the payee order, and the operator's response to the payee order is encrypted with the session key contained in the payee order. The payee judges the performing of the payment according to the validity of the signature on the payee receipt.

In the payee device, by means of the operator's response to the payee order, one forms and delivers to the payer device data confirming the payee's consent to the fact that the payment operation has been successfully performed.

The operator opens a payment account connected with each of the payment certificates and connects it with the public key of the payment certificate, and the funds of the payment account are spent in one or several payment operations. Here, the public key account whose public key coincides with the public key of the payment certificate is used as the payment account. The funds appear on the payment account as a result of one or several operations of crediting the payment account performed from the funds of the value contained in the payment certificate, and the operation of opening a payment account is combined with the operation of crediting this account that was first in time.

When performing the operation of crediting a payment account, one delivers to the payment server a payment certificate signature, whose level can be chosen arbitrarily within the level of the payment certificate. Moreover, each operation of crediting a payment account is combined with one of the payment operations.

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The possibility of realization of the above-described best embodiment of the method for effecting payments by the third variant is clarified by the following example.

Example 4

In this example, we use agreements and notation adopted in Example 1. The preparatory actions, opening the seller's account, replenishment of the payer device, and the seller's actions during the execution of a payment operation are performed in this example as in Example 2 described above.

Examples of three payment operations are given below. Here, the first one of the payment operations described below, which is the first (in time) operation performed with the use of a payment certificate, is combined with the operation of opening a payment account and its crediting. The second one of the payment operations described below is combined with the operation of crediting an already opened payment account. The third one of the payment operations described below is performed under the conditions when the funds on the payment account already suffice for performing such a payment operation.

Below, given is a description of the first (in time) payment operation performed with the use of a payment certificate, which is combined with the operation of opening a payment account and its crediting.

Suppose that the payer device contains a payment certificate with base (EP, X) and signature S of level $(2, 12, 45)$. The payment certificate value is equal to 212 dollars and 45 cents. Suppose, furthermore, that the payment certificate has not already been used in one of the payment operations.

The payer willing to pay the seller an amount of 18.999 dollars for some goods, prepares the payment data *PaymentData*, including in them the payer order *PayerOrder* signed with the secret key of the payment certificate DP , and the data A intended for the seller and consisting in this example of the name of the paid goods and the identification data of the recipient of the goods.

The payer order *PayerOrder* consists of the public key EP of the payment certificate, the signature S_1 of the payment certificate, the seller's account number *AccountId*, and the data C defining the conditions of payment. In the present example, the payer takes as C the seller's account number *AccountId* and the value of the hash function H on the text of the seller's obligation valid in the case of a payment, 3039 namely, an obligation to supply the corresponding goods to the person with the indicated identification data.

The public key EP included in the payer order *PayerOrder* is used in the payment server for opening the payment account connected with this key. The signature S_1 of the payment certificate included in the payer order *PayerOrder* is used in the payment server for crediting the payment account with public key EP .

The signature S_1 is made in the payer device from the signature S of the payment certificate. Here, the level $L=(L_1, \dots, L_3)$ of the signature S_1 is chosen in so a way that, on the one hand, the value corresponding to this level would suffice for performing the present payment operation, and, on the other hand, it would not exceed the level of the signature S in each of the three orders. Another aim of such a choice is to hide the level of the signature S contained in the payer device from the operator. In the present example, the level of the signature S is represented by the collection $(2, 12, 45)$, and the value sufficient for performing the present payment operation, expressed in cents, is, in the present case, 1899.9

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cents. Thus, the restrictions on the choice of the level $L=(L_1, \dots, L_3)$ are as follows: first, L_1 does not exceed 2, L_2 does not exceed 12, L_3 does not exceed 45, and, secondly, the quantity $L_1 \cdot N_1 + L_2 \cdot N_2 + L_3 \cdot N_3$, where $N_1 = 100^2$, $N_2 = 100$, $N_3 = 1$, must be at least 1899.9 cents. The level L satisfying these conditions is chosen with the help of a random number generator and, in the present example, $L = (1, 3, 14)$. The signature S_I is made with the help of a computing device programmed in the corresponding way; in accordance with the relation $S_I = S^V \pmod{N}$, where $V = V_1 \cdot V_2 \cdot V_3$, $V_1 = E_1^{M_1-L_1}$, $V_2 = E_2^{M_2-L_2}$, $V_3 = E_3^{M_3-L_3}$, and the level (M_1, \dots, M_3) is equal to the level $(2, 12, 45)$, i.e., to the level of the signature S .

The seller willing to accept the payment forms his payee order *SellerOrder* = (*AccountId*, *PayerOrder*) encrypted with the bank's public encryption key, and delivers it to the bank.

Having made sure that the payment account database contains no record of a payment account with public key EP and having verified the validity of the signature S_I with the help of a public or a secret money key, the bank opens a payment account and enters the corresponding record into the payment account database. Here, first, the opened account is credited with the amount corresponding to the level $L = (1, 3, 14)$ of the signature S_I , i.e., in the present example, with an amount of 103 dollars and 14 cents. In addition, one debits the opened payment account with an amount of 50 cents, which is equal to the charge for opening the payment account. After that, the payment operation per se is performed. Namely, having verified the signature on the payer order *PayerOrder* with the help of the public key EP of the payment account, the bank enters the signed payer order *PayerOrder* into the information storage, debits the payment account with an amount of 1899.9 cents and credits the account with number *AccountId* with an amount of 1799.9 cents, under the assumption that the cost of performing a payment operation by the bank is equal to 1 cent. After that, the bank forms the payer receipt confirming the fact of crediting the account with number *AccountId* with an amount of 1799.9 cents, signs it, and delivers to the seller, who, having verified the validity of the bank's signature on the receipt received, regards the payment as being effected and informs the payer about successfully effecting the payment.

A description of a payment operation, which is combined with the operation of crediting an already opened payment account is given below. Generally speaking, the payee participating in this operation is in no way connected with the payee from the first payment operation described above.

Suppose that the payer device contains a payment certificate with base (EP, X) and signature S of level $(2, 12, 45)$. The payment certificate value is equal to 212 dollars and 45 cents. Suppose that the payment account connected with the given certificate is already opened, the total amount of expenses of this payment account does not exceed the amount of 6732.8 cents, and the level of the signature delivered to the operator in one of the previous operations of crediting the payment account does not exceed the level $(1, 3, 14)$. Thus, the payment account has been credited with an amount of 10 314 cents in the course of the previous crediting operations.

The payer willing to pay the seller an amount of 3699.9 cents for some goods prepares the payment data *PaymentData*, including in them the payer order *PayerOrder* signed with the secret key DP of the payment certificate, and the data A intended for the seller and consisting in this example of the name of the paid goods and the identification data of the re-

recipient of the goods.

Since the payment amount exceeds the difference between the amount of 10 314 cents, delivered to the payment account in the previous crediting operations, and the total amount of expenses of this payment account, which makes 6732.8 cents, the payer combines the payment operation with the operation of crediting a payment account.

The payer order *PayerOrder* consists of the identifier of the public key of the payment certificate, which coincides in the present example with the value on the key *EP* of the hash function *H* fixed in advance, the signature S_I of the payment certificate, the seller's account number *AccountId*, and the data *C*, defining the conditions of payment. In the present example, the payer takes as *C* the seller's account number *AccountId* and the value of the hash function *H* on the text of the seller's obligation valid in the case of a payment, namely, an obligation to supply the corresponding goods to the person with the indicated identification data.

The identifier $H(EP)$ of the public key of the payment certificate included in the payer order *PayerOrder* is used in the payment server for search of the payment account connected with the key *EP*. The signature S_I of the payment certificate included in the payer order *PayerOrder* is used in the payment server for crediting the payment account with public key *EP*.

The signature S_I is made in the payer device from the signature *S* of the payment certificate. Here, the level $L=(L_1, \dots, L_3)$ of the signature S_I is chosen so that, first, the value corresponding to this level suffices for performing the present payment operation considering the expenses paid earlier, secondly, that said level does not exceed the level of the signature *S* in each of the three orders, and, thirdly, that in each of the three orders the said level would not be less than the level of the signature delivered to the operator in one of the previous operations of crediting the payment account, i.e., in the present example, not less than the level (1, 3, 14). In the present example, one takes $L = (2, 7, 15)$. The signature S_I is made with the help of a computing device programmed in the corresponding way, in accordance with the relation $S_I = S^V \pmod{N}$, where $V = V_1 \cdot V_2 \cdot V_3$, $V_1 = E_1^{M1-L1}$, $V_2 = E_2^{M2-L2}$, $V_3 = E_3^{M3-L3}$, and the level (M_1, \dots, M_3) is equal to the level (2, 12, 45), i.e., to the level of the signature *S*.

The seller willing to accept the payment forms the payee order *SellerOrder* = (*AccountId*, *PayerOrder*) encrypted with the bank's public encryption key and delivers it to the bank.

The bank, having detected a record of a payment account with public key *EP* in the payment account database and having verified the validity of the signature S_I with the help of a public or a secret money key, increases the total amount on the found account to the amount corresponding to the level each order of which equals the maximal one of the corresponding orders of the levels of all signatures of the payment certificates delivered to the bank in the operations of crediting the given payment account. After that, the payment operation per se is performed. This operation is performed as in the above example of the first payment operation.

Below, a description of a payment operation performed with the help of an already opened payment account without combining it with the operation of crediting a payment account is given. Generally speaking, the payee participating in this operation is in no way connected with the payee from the first payment operation described above.

Such an operation is performed in the case where the payment amount does not exceed

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the difference between the amount delivered to the payment account in previous crediting operations and the total amount of expenses of this payment account. In such an operation, the payer order *PayerOrder* does not contain the payment certificate signature, but consists of the identifier of the public key of the payment certificate, which identifier coincides in the present example with the value on the key *EP* of the hash function *H* fixed in advance, the seller's account number *AccountId*, and the data *C* determining the conditions of payment. In the present example, the payer takes as *C* the seller's account number *AccountId* and the value of the hash function *H* on the text of the seller's obligation valid in the case of a payment, namely, an obligation to supply the corresponding goods to the person with the indicated identification data. In all other respects, such an operation proceeds in the same way as the payment operations described above.

A description of the best embodiment of the method for effecting payments by the fourth variant is given below.

In the best embodiment of the method for effecting payments by the fourth variant, one performs all actions that are performed in the best embodiment of the method for effecting payments by the third variant. In addition, one additionally performs a payment certificate replenishment operation, and the best embodiment of this operation is described above in the description of the best embodiment of the method for effecting payments by the second variant.

The possibility of realization of the above-described best embodiment of the method for effecting payments by the fourth variant is clarified by the above-described Example 4, where the possibility of realization of the best embodiment of the method for effecting payments by the third variant is clarified, and by the above-described Example 3, where the possibility of realization of the best embodiment of replenishing the payer device with the help of replenishing the payment certificate is clarified.

A description of the best embodiment of the apparatus for effecting payments is given below.

In the best embodiment, the payer device contains a means for replenishing the payer device by the use of making a blind money signature of an operator, which is realized by a means for increasing the level of the payment certificate signature. In addition, the payer device contains a means for opening a public key account. Furthermore, the means for forming a payer order signed with a secret key of the payment certificate has a means for forming a demand for crediting the payment account, which, in turn, has a means for decreasing the level of the payment certificate signature.

The payee device contains a means for opening a public key account and a means for verifying the signed payee receipt with the operator's public key.

The payment server contains a means for serving a database of payment accounts and a means for serving a database of accounts, and the means for serving a database of accounts has a means for opening a public key account, a means for crediting an account and a means for debiting an account. The means for serving a database of payment accounts has a means for opening a payment account, a means for verifying the money signature and a means for crediting a payment account. The means for performing a payment operation contained in the payment server has a means for verifying a signature on the payer order and a means for making a signed payee receipt. In addition, the payment server contains a means for making

The means for storing information in the storage devices of the payer device, of the payee device, and of the payment server have high reliability, and the payer device, the payee device, and the payment server are equipped with means for encryption of outgoing messages and means for decryption of incoming messages.

Example 5

Fig.2 shows a flow-chart of a payer device replenishment operation. Here, block 4 shows a means for creating a payment certificate base by processing a public key of a payment certificate with a one-way function, block 5 shows a storage device, line 6 shows a means for storing the created payment certificate base in the storage device, block 7 shows a means for increasing the level of the payment certificate signature, block 8 shows a means for forming a money demand including a blinded payment certificate signature, block 9 shows a means for unblinding the data to be unblinded comprised in a response to the money demand, line 10 shows a means for entering the result of unblinding into the storage device, block 11 shows a means for processing the money demand, and block 12 shows a means for making a money signature. In addition, line 13 shows a means for reading the payment certificate signature from the storage device, line 14 shows a connection via which the money demand is delivered to the payment server, and line 15 shows a connection via which the operator's response to the money demand is delivered to the payer device.

The apparatus for effecting a payment is realized by programming means on the basis of the algorithms indicated in Examples 2 and 3 above. In particular, cryptographic means used, such as making a signature, verification of a signature, encryption, and decryption, are based on the functions of arithmetic of integral numbers and modular arithmetic. Examples of realization of such functions are well known in the art. Means for calculating hash functions utilized are also well known.

With the help of the apparatus for effecting payments, which is shown in Fig.1, payments are effected as follows.

Payer device 2 is replenished for an arbitrary number of times. Here, one can replenish the payer device with the help of both primarily filling a payment certificate and replenishing one of the payment certificates already at hand. When replenishing the payer device with the help of primarily filling a payment certificate, one uses the means for creating a payment certificate base, in which the public key of the payment certificate is processed by a one-way function, after which the created payment certificate base is entered via line 6 into storage device 5, both as the base of the payment certificate and as the payment certificate signature of zero level.

In an arbitrary operation of replenishing payer device 2, a payment certificate signature is entered via line 13 into block 7, where the payment certificate signature is blinded and a money demand including the blinded payment certificate signature is formed with the help of block 8. The money demand is delivered via line 14 to payment server 1, where it is processed in block 11, including making a money signature on the blinded payment certificate signature with the help of block 12. The operator's response to the money demand, which response is formed in block 11, is entered via line 15 into payer device 2, where the data to be unblinded contained in the response to the money demand is unblinded with the help of block 9. Unblinded data are entered via line 10 into storage device 6 as the payment certificate signature of higher level.

When performing a payment operation, a payer order signed with the secret key of the payment certificate is formed in block 16 of payer device 2. The payer order is entered via line 22 into payee device 3, where a payee order including the payer order is formed in block 17. The payee order is entered via line 23 into payment server 1, where with the help of block 19 one performs a payment operation, in the course of which the payee's account stored in account database 21 is credited from the funds of the payment account stored in payment account database 20. Here, reading and modification of records of the payment account and the payee's account are performed via lines 24 and 25, respectively. The operator's response to the payee order formed in block 19 and including a payee receipt signed with the operator's secret key is entered via line 26 into block 18 of payee device 3, where the signed payee receipt is verified, which completes the payment operation.

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Variants of implementing the invention

For each of the variants of the claimed method for effecting payments, the invention admits an embodiment such that the payer device can be replenished from the funds of an intermediary payer. In this case, when replenishing his payer device, the payer delivers the blinded data to an intermediary payer, who additionally blinds said blinded data and, correspondingly, unblinds the data to be unblinded which he receives from the bank.

If need be, the payee can also control the incoming payments. Such control is intended in order to ensure that nobody can credit the payee's account without his assistance. To exercise such control, the payee order and its signature with the payee's secret key are delivered to the payment server, and during the performing of a payment operation the validity of the signature on the payee order is verified. Furthermore, when forming the payer and payee orders, the conditions of payment are included in these orders, and during the execution of a payment operation the correspondence between the conditions of payment contained in the payer and payee orders is controlled in the payment server. In particular, in the payer device

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when forming the payment data and in the payee device when forming the payee order, one processes the payee obligation data with the help of one and the same one-way function, and the data obtained by the processing are included both in the payment data and in the payee order as a part of the conditions of payment.

- 5 The reliability of the method for effecting payments by each of the variants is ensured, in particular, by that if there are malfunctions in communication nets when performing the operations employed in a payment, such operations can be recovered till their successful completion without damage for the parties involved.

- 10 As the replenishment source in the payer device replenishment operation, one can use the payer's account, which was credited in advance when performing an earlier payment operation, in which the payer, in turn, played a role of the payee.

When performing a payment operation, the payee can play a role of a payer. Such a payment operation can be used for transferring values from the payer device to the payer's account.

- 15 An overgrowth of the operator's databases storing information on payment certificates can be prevented by charging either the operation of replenishing the payer device or the operation of opening the payment account connected with a payment certificate.

- 20 The operator's money obligations associated with payment certificates can be expressed in various currencies, and performing both a payment operation and the payer device replenishment operation can be combined with the operation of converting one currency into another.

- 25 To increase the security of the participants of a payment system utilizing the present invention, one can introduce certain restrictions both on the amount of a single replenishment of the payer device, and on the total amount of expenses of the replenishment source within a certain period of time.

The message exchange between the payee device and the payment server, the payer device and the payment server, and the payee device and the payer device can proceed in the interactive mode. In particular, the money demand, the payment data, the payee order, and other data can be delivered to their addressee in portions.

- 30 The procedure of acknowledging by the payment system operator of his obligations associated with a payment certificate, besides the verification of the operator's own signature of the certificate, can include verification of the validity period of other data, as well.

- 35 In addition, the data to be unblinded, which one received when replenishing the payer device, together with data allowing one to perform such an unblinding can be used as the payment certificate signature, because they allow one to convince a third party in the existence of the operator's obligations.

Industrial applicability

- 40 The invention can be used in electronic queueing systems, especially those where payments via open communication nets are required. Among possible applications, the invention can be used for organizing payment systems, trade systems, service centers, and in many other areas. In particular, the invention can be used in the work of banks and systems of banks, for organizing shops, trade in securities, lotteries, etc.

Claims

1. A method for effecting payments comprising steps of: replenishing a payer device by an operation of primarily filling a payment certificate, in which a payment certificate base is created in the payer device and a payment certificate signature is obtained by means of making a blind money signature of an operator; performing an operation of opening a payee's account; performing a payment operation in which the payment certificate signature and an identifier of the payment certificate base are included into payment data delivered to a payee device by means of which a payee order is formed including the payment certificate signature and the identifier of the payment certificate base; delivering the formed payee order to a payment server in which the payee's account is credited on the basis of the payee order in the case of absence of an information that the payment certificate was utilized, according to the validity of the delivered payment certificate signature; and forming the operator's response to the payee order, by means of which response the payment operation is judged, *characterized* in that the method is further comprising steps of: including an identifier of a public key into the payment certificate base, said public key corresponding to an arbitrary secret key of a payer, wherein the public key is accepted as a public key of the payment certificate and the secret key of the payer is accepted as a secret key of the payment certificate; including a payer order signed with the secret key of the payment certificate into the payment data, and including an information on the payee and the identifier of the payment certificate base into the payer order; and the step of crediting the payee's account is carried out according to the validity of the signature on the payer order.
2. The method according to claim 1, *characterized* in that in the step of performing the payment operation the signed payer order is entered into the operator's information storage.
3. The method according to claim 1, *characterized* in that in the step of replenishing the payer device a money demand is formed including data for making a blind money signature, and is delivered to the payment server in which a replenishment source and replenishment amount are determined according to the money demand; data to be unblinded are created in the step of making the blind money signature by processing data for making the blind money signature, that are comprised in the demand, with a secret money key corresponding to the replenishment amount, whereupon the payment certificate signature is made in the payer device by unblinding.
4. The method according to claim 3, *characterized* in that in the step of replenishing the payer device by the operation of primarily filling the payment certificate a blinded identifier of the created payment certificate base is included as the data for making the blind money signature into the money demand being formed.
5. The method according to claim 1, *characterized* in that in the step of performing the payment operation a payee receipt is included into the operator's response to the payee order, which receipt being signed with the arbitrary secret key of the operator, and the performing of the payment for the payee is judged according to the validity of the signature on the payee receipt.
6. The method according to claim 1, *characterized* in that in the step of performing the payment operation data are formed in the payee device with the use of the operator's response to the payee order and delivered to the payer device, according to which data the

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7. The method according to claim 6, *characterized* in that in the step of performing the payment operation a payer receipt is included into the operator's response to the payee order and into the data delivered to the payer device, which receipt being signed with the arbitrary secret key of the operator, and the performing of the payment for the payer is judged according to the validity of the signature on the payer receipt.

10 9. The method according to claim 1, *characterized* in that in performing operations with the payment certificate the public key of the payment certificate converted by an arbitrary one-way function is used as the identifier of the payment certificate base.

15 the public key of the payment certificate.
11. The method according to claim 1, *characterized* in that in the step of replenishing the payer device the validity of the made payment certificate signature is verified.

12. The method according to claim 1, *characterized* in that in the step of opening the account an arbitrary secret key is accepted as the secret key of the account, and the public key corresponding to the secret key of the account is delivered to the payment server as a public key of the account being opened.

14. The method according to claim 13, *characterized* in that payee obligation data are in-
25 cluded into the conditions of payment comprised in the payer order.

30 16. The method according to claim 13, *characterized* in that in the payer device in the step of forming the payment data, the payee obligation data are processed by an arbitrary one-way function, and data obtained in this processing are included into the payer order as the conditions of payment.

18. The method according to claim 1, *characterized* in that in the step of replenishing the payer device a payer's account is used as a replenishment source.

20. The method according to claim 1, *characterized* in that in the step of performing the payment operation the payee appears as the payer.

21. The method according to claim 1, *characterized* in that in the step of performing the payment operation a part of a payment certificate value is returned to the payer device.

22. The method according to claim 1, *characterized* in that the step of replenishing the payer device is performed from funds of an intermediate payer.

23. The method according to claim 22, *characterized* in that in the step of replenishing the payer device, data blinded in the payer device in the step of making the blind money signature of the operator are subjected to an additional blinding in the payer device of the intermediate payer.

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24. A method for effecting payments comprising steps of: replenishing a payer device by an operation of primarily filling a payment certificate, in which a payment certificate base is created in the payer device and a payment certificate signature is obtained by means of making a blind money signature of an operator; performing an operation of opening a payee's account; performing a payment operation in which the payment certificate signature and an identifier of the payment certificate base are included into payment data delivered to a payee device by means of which a payee order is formed including the payment certificate signature and the identifier of the payment certificate base; delivering the formed payee order to a payment server in which the payee's account is credited on the basis of the payee order in the case of absence of an information that the payment certificate was utilized, according to the validity of the delivered payment certificate signature; and forming the operator's response to the payee order, by means of which response the payment operation is judged, *characterized* in that the method is further comprising steps of: including an identifier of a public key into the payment certificate base, said public key corresponding to an arbitrary secret key of the payee, wherein the public key is accepted as a public key of the payment certificate and the secret key of the payee is accepted as a secret key of the payment certificate; performing the replenishment of the payer device by means of an operation of replenishing the payment certificate, in which operation a blind money signature of the operator on the payment certificate signature already being in the payer device is made; including a payer order signed with the secret key of the payment certificate into the payment data, and including an information on the payee and the identifier of the payment certificate base into the payer order; and the step of crediting the payee's account is carried out according to the validity of the signature on the payer order.
25. The method according to claim 24, *characterized* in that in the step of replenishing the payer device a money demand is formed including data for making a blind money signature, and is delivered to the payment server in which a replenishment source and replenishment amount are determined according to the money demand; data to be unblinded are created in the step of making the blind money signature by processing data for making the blind money signature, that are comprised in the demand, with a secret money key corresponding to the replenishment amount, whereupon the payment certificate signature is made in the payer device by unblinding.
26. The method according to claim 25, *characterized* in that in the step of replenishing the payer device by the operation of primarily filling the payment certificate a blinded identifier of the created payment certificate base is included as the data for making the blind money signature into the money demand being formed.
27. The method according to claim 25, *characterized* in that in the step of replenishing the payer device by the operation of replenishing the payment certificate a blinded payment certificate signature is included as the data for making the blind money signature into the money demand being formed.
28. The method according to claim 24, *characterized* in that in the step of performing the payment operation the signed payer order is entered into the operator's information storage.
29. The method according to claim 24, *characterized* in that in the step of performing the payment operation a payee receipt is included into the operator's response to the payee order, which receipt being signed with the arbitrary secret key of the operator, and the per-

forming of the payment for the payee is judged according to the validity of the signature on the payee receipt.

30. The method according to claim 24, *characterized* in that in the step of performing the payment operation data are formed in the payee device with the use of the operator's response to the payee order and delivered to the payer device, according to which data the performing of the payment for the payer is judged.

31. The method according to claim 30, *characterized* in that in the step of performing the payment operation a payer receipt is included into the operator's response to the payee order and into the data delivered to the payer device, which receipt being signed with the arbitrary secret key of the operator, and the performing of the payment for the payer is judged according to the validity of the signature on the payer receipt.

32. The method according to claim 31, *characterized* in that the payer receipt is encrypted by an arbitrary encryption key of the payer prior to including said receipt into the operator's response to the payee order.

33. The method according to claim 24, *characterized* in that in performing operations with the payment certificate the public key of the payment certificate converted by an arbitrary one-way function is used as the identifier of the payment certificate base.

34. The method according to claim 24, *characterized* in that in performing operations with the payment certificate the public key of the payment certificate is used as the identifier of the public key of the payment certificate.

35. The method according to claim 24, *characterized* in that in the step of replenishing the payer device the validity of the made payment certificate signature is verified.

36. The method according to claim 24, *characterized* in that in the step of opening the account an arbitrary secret key is accepted as the secret key of the account, and the public key corresponding to the secret key of the account is delivered to the payment server as a public key of the account being opened.

37. The method according to claim 24, *characterized* in that conditions of payment are included into the payer order.

38. The method according to claim 37, *characterized* in that payee obligation data are included into the conditions of payment comprised in the payer order.

39. The method according to claim 38, *characterized* in that prior to performing the payment operation, the payee obligation data are signed with an arbitrary secret key of the payee, and the payer verifies the payee's signature on the payee obligation data prior to performing the payment operation.

40. The method according to claim 38, *characterized* in that in the payer device in the step of forming the payment data, the payee obligation data are processed by an arbitrary one-way function, and data obtained in this processing are included into the payer order as the conditions of payment.

41. The method according to claim 24, *characterized* in that the payer order is encrypted by an arbitrary public encryption key of the operator prior to including them into the payment data.

42. The method according to claim 24, *characterized* in that in the step of replenishing the payer device a payer's account is used as a replenishment source.

43. The method according to claim 24, *characterized* in that in the step of replenishing the

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payer device a bank card is used as a replenishment source.

44. The method according to claim 24, *characterized* in that in the step of performing the payment operation the payee appears as the payer.

5 45. The method according to claim 24, *characterized* in that in the step of performing the payment operation a part of a payment certificate value is returned to the payer device.

46. The method according to claim 24, *characterized* in that the step of replenishing the payer device is performed from funds of an intermediate payer.

10 47. The method according to claim 46, *characterized* in that in the step of replenishing the payer device, data blinded in the payer device in the step of making the blind money signature of the operator are subjected to an additional blinding in the payer device of the intermediate payer.

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54. The method according to claim 48, *characterized* in that in the step of performing the

payment operation a payee receipt is included into the operator's response to the payee order, which receipt being signed with the arbitrary secret key of the operator, and the performing of the payment for the payee is judged according to the validity of the signature on the payee receipt.

- 5 55. The method according to claim 48, *characterized* in that in the step of performing the payment operation data are formed in the payee device with the use of the operator's response to the payee order and delivered to the payer device, according to which data the performing of the payment for the payer is judged.
56. The method according to claim 55, *characterized* in that in the step of performing the
10 payment operation a payer receipt is included into the operator's response to the payee order and into the data delivered to the payer device, which receipt being signed with the arbitrary secret key of the operator, and the performing of the payment for the payer is judged according to the validity of the signature on the payer receipt.
57. The method according to claim 56, *characterized* in that the payer receipt is encrypted
15 by an arbitrary encryption key of the payer prior to including said receipt into the operator's response to the payee order.
58. The method according to claim 48, *characterized* in that in performing operations with the payment certificate the public key of the payment certificate converted by an arbitrary one-way function is used as the identifier of the payment certificate base.
- 20 59. The method according to claim 48, *characterized* in that in performing operations with the payment certificate the public key of the payment certificate is used as the identifier of the public key of the payment certificate.
60. The method according to claim 48, *characterized* in that in the step of replenishing the payer device the validity of the made payment certificate signature is verified.
- 25 61. The method according to claim 48, *characterized* in that in the step of opening the account an arbitrary secret key is accepted as the secret key of the account, and the public key corresponding to the secret key of the account is delivered to the payment server as a public key of the account being opened.
62. The method according to claim 48, *characterized* in that conditions of payment are included
30 into the payer order.
63. The method according to claim 62, *characterized* in that payee obligation data are included into the conditions of payment comprised in the payer order.
64. The method according to claim 63, *characterized* in that prior to performing the payment operation, the payee obligation data are signed with an arbitrary secret key of the
35 payee, and the payer verifies the payee's signature on the payee obligation data prior to performing the payment operation.
65. The method according to claim 62, *characterized* in that in the payer device in the step of forming the payment data, the payee obligation data are processed by an arbitrary one-way function, and data obtained in this processing are included into the payer order as the
40 conditions of payment.
66. The method according to claim 48, *characterized* in that the payer order is encrypted by an arbitrary public encryption key of the operator prior to including them into the payment data.
67. The method according to claim 48, *characterized* in that in the step of replenishing the

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payer device a payer's account is used as a replenishment source.

68. The method according to claim 48, *characterized* in that in the step of replenishing the payer device a bank card is used as a replenishment source.

5 69. The method according to claim 48, *characterized* in that in the step of performing the payment operation the payee appears as the payer.

70. The method according to claim 48, *characterized* in that in the step of performing the payment operation a part of a payment certificate value is returned to the payer device.

10 71. The method according to claim 48, *characterized* in that in the step of replenishing the payer device, a payment account associated with the base of one of payment certificates is used as a replenishment source.

72. The method according to claim 48, *characterized* in that the step of replenishing the payer device is performed from funds of an intermediate payer.

15 73. The method according to claim 72, *characterized* in that in the step of replenishing the payer device, data blinded in the payer device in the step of making the blind money signature of the operator are subjected to an additional blinding in the payer device of the intermediate payer.

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74. A method for effecting payments comprising steps of: replenishing a payer device by an operation of primarily filling a payment certificate, in which a payment certificate base is created in the payer device and a payment certificate signature is obtained by means of making a blind money signature of an operator; performing an operation of opening a payee's account; performing a payment operation in which an identifier of the payment certificate base is included into payment data delivered to a payee device by means of which a payee order is formed including the identifier of the payment certificate base received from a payer; delivering the formed payee order to a payment server in which the payee's account is credited on the basis of the payee order; forming the operator's response to the payee order, by means of which response the payment operation is judged, *characterized* in that the method is further comprising steps of: including an identifier of a public key into the payment certificate base, said public key corresponding to an arbitrary secret key of a payer, wherein the public key is accepted as a public key of the payment certificate and the secret key of the payer is accepted as a secret key of the payment certificate; performing the operation of replenishing the payer device by an operation of replenishing the payment certificate, in which a blind money signature of the operator on the payment certificate signature already being in the payer device is made; opening a payment account associated with the payment certificate base; carrying out the operation of crediting the payment account, in which a payment certificate signature is delivered to the payment server, choosing the level of said signature arbitrarily within the level of the payment certificate, and the operation of crediting the payment account is carried out in accordance with the excess of the level of the delivered signature above the level of the payment account; including a payer order signed with the secret key of the payment certificate into the payment data, and including an information on the payee and the identifier of the payment certificate base into the payer order; and the step of crediting the payee's account is carried out according to the validity of the signature on the payer order from funds of the payment account.

75. The method according to claim 74, *characterized* in that the payment account associated with the payment certificate base is opened in the step of performing the payment operation.

76. The method according to claim 74, *characterized* in that the crediting operation is carried out in the step of performing the payment operation.

77. The method according to claim 74, *characterized* in that in the step of performing the payment operation the signed payer order is entered into the operator's information storage.

78. The method according to claim 74, *characterized* in that in the step of replenishing the payer device a money demand is formed including data for making a blind money signature, and is delivered to the payment server in which a replenishment source and replenishment amount are determined according to the money demand; data to be unblinded are created in the step of making the blind money signature by processing data for making the blind money signature, that are comprised in the demand, with a secret money key corresponding to the replenishment amount, whereupon the payment certificate signature is made in the payer device by unblinding.

79. The method according to claim 78, *characterized* in that in the step of replenishing the payer device by the operation of primarily filling the payment certificate a blinded identifier

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80. The method according to claim 78, *characterized* in that in the step of replenishing the payer device by the operation of replenishing the payment certificate a blinded payment certificate signature is included as the data for making the blind money signature into the money demand being formed.

82. The method according to claim 74, *characterized* in that in the step of performing the payment operation data are formed in the payee device with the use of the operator's response to the payee order and delivered to the payer device, according to which data the performing of the payment for the payer is judged.

84. The method according to claim 83, *characterized* in that the payer receipt is encrypted by an arbitrary encryption key of the payer prior to including said receipt into the operator's response to the payee order.

86. The method according to claim 74, *characterized* in that in performing operations with the payment certificate the public key of the payment certificate is used as the identifier of the public key of the payment certificate.

88. The method according to claim 74, *characterized* in that in the step of opening the account an arbitrary secret key is accepted as the secret key of the account, and the public key corresponding to the secret key of the account is delivered to the payment server as a public key of the account being opened.

90. The method according to claim 89, *characterized* in that payee obligation data are included into the conditions of payment comprised in the payer order.

92. The method according to claim 89, *characterized* in that in the payer device in the step

of forming the payment data, the payee obligation data are processed by an arbitrary one-way function, and data obtained in this processing are included into the payer order as the conditions of payment.

5 93. The method according to claim 74, *characterized* in that the payer order is encrypted by an arbitrary public encryption key of the operator prior to including them into the payment data.

94. The method according to claim 74, *characterized* in that in the step of replenishing the payer device a payer's account is used as a replenishment source.

10 95. The method according to claim 74, *characterized* in that in the step of replenishing the payer device a bank card is used as a replenishment source.

96. The method according to claim 74, *characterized* in that in the step of performing the payment operation the payee appears as the payer.

97. The method according to claim 74, *characterized* in that in the step of performing the payment operation a part of a payment certificate value is returned to the payer device.

15 98. The method according to claim 74, *characterized* in that in the step of replenishing the payer device, a payment account associated with the base of one of payment certificates is used as a replenishment source.

99. The method according to claim 74, *characterized* in that the step of replenishing the payer device is performed from funds of an intermediate payer.

20 100. The method according to claim 99, *characterized* in that in the step of replenishing the payer device, data blinded in the payer device in the step of making the blind money signature of the operator are subjected to an additional blinding in the payer device of the intermediate payer.

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101. An apparatus for effecting payments comprising a payer device, payee device and payment server interconnected by telecommunication nets, the payer device comprising a means for replenishing the payer device by the use of making a blind money signature of an operator, and the payment server comprising a means for making a money signature,
 5 *characterized* in that the payer device further comprises a means for creating a payment certificate base by processing a public key of the payment certificate with a one-way function, a means for storing the created payment certificate base in a storage device, and a means for forming a payer order signed with a secret key of the payment certificate; the payee device comprises a means for forming a payee order including the payer order; the
 10 payment server further comprises a means for performing a payment operation, a means for serving a database of payment accounts, and a means for serving a database of accounts, wherein said means for performing a payment operation has a means for verifying a signature on the payer order and a means for making a signed payee receipt, said means for serving the database of payment accounts has a means for verifying the money signature, and the means for replenishing the payer device by the use of making the blind money
 15 signature of the operator is realized using a means for increasing the level of the payment certificate signature.

102. The apparatus according to claim 101, *characterized* in that the payee device comprises a means for opening a public key account, and said means for serving the database
 20 of accounts has a means for opening a public key account.

103. The apparatus according to claim 101, *characterized* in that the payer device comprises a means for opening a public key account, and said means for serving the database of accounts has a means for opening a public key account.

104. The apparatus according to claim 101, *characterized* in that said means for increasing the level of the payment certificate signature has a means for forming a money demand including a blinded payment certificate signature, a means for unblinding the data to be
 25 unblinded comprised in a response to the money demand, and a means for entering the result of unblinding into said storage device, and the payment server comprises a means for processing the money demand, wherein said means for processing the money demand has a means for making the money signature.
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105. The apparatus according to claim 101, *characterized* in that said means for serving the database of payment accounts has a means for opening a payment account and a means for crediting a payment account.

106. The apparatus according to claim 101, *characterized* in that the payee device has a means for verifying the signed payee receipt.
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107. The apparatus according to claim 101, *characterized* in that said means for forming the payer order signed with the secret key of the payment certificate has a means for forming a demand for crediting the payment account.

108. The apparatus according to claim 101, *characterized* in that said means for forming the demand for crediting the payment account has a means for decreasing the level of the payment certificate signature.
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109. The apparatus according to claim 101, *characterized* in that the payer device, payee device and payment server are further provided with a means for encryption of outgoing messages and a means for decryption of incoming messages.

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Abstract

- A method for effecting payments is suggested which allows protecting the financial interests of each participant of a payment during payments via open telecommunication nets from other participants' cheating, provides protection of the payers' and payees' privacy, admits payments ranging from micro-payments to business-to-business payments, ensures that the time taken to effecting a payment depends only on the speed of action of the net connections and not on the payment amount, makes it possible to serve of serving a number of clients which grows proportionally to the payment system operator's resources, is easy to build into an arbitrary trade system, enables each client both to pay and to receive payments, and makes possible payments between clients of different banks. The method of effecting payments is realized with the help of programming means.

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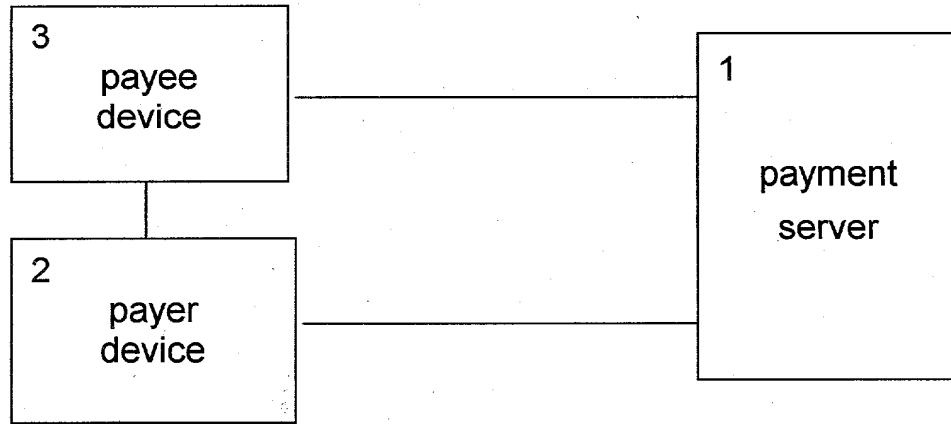


Fig. 1

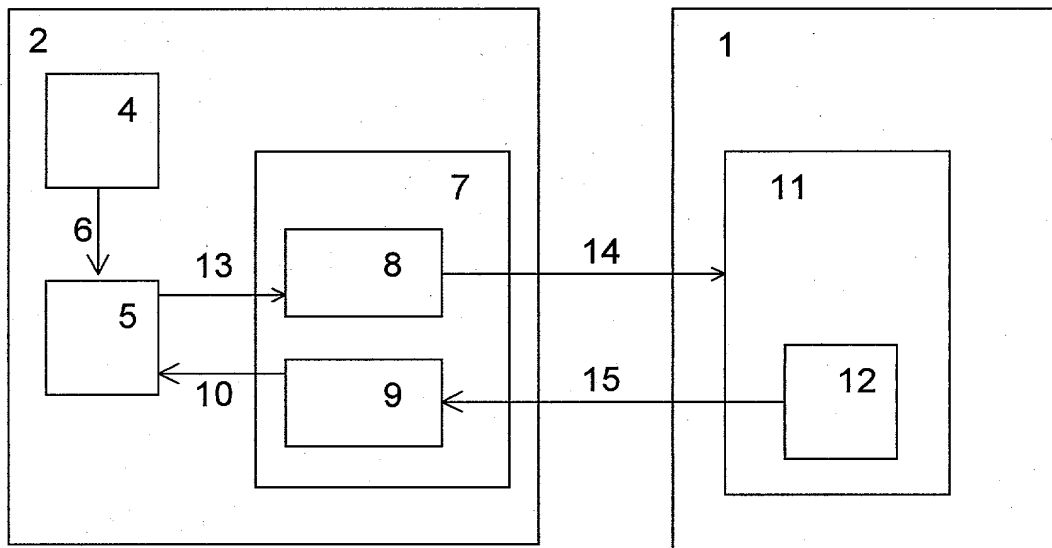


Fig. 2

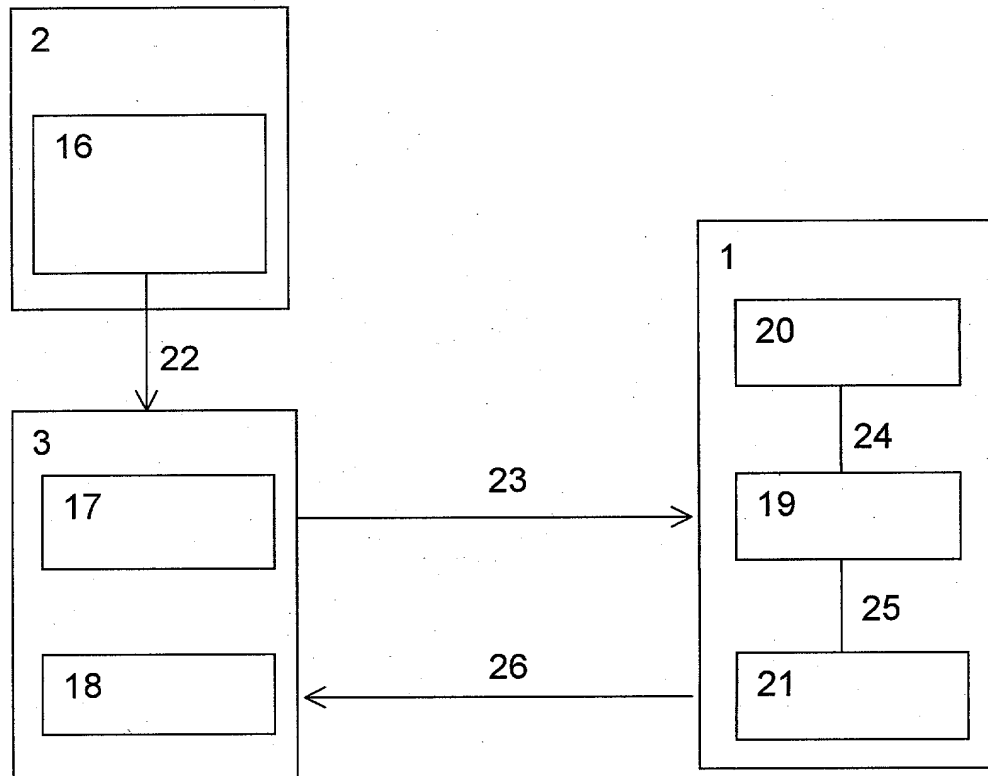


Fig. 3

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Attorney Docket Number: P-9902 MG

First Named Inventor: Zolotarev

Application No.: (not yet known)

Filing Date: Group Art Unit: Examiner Name: (not yet known)

UNITED STATES - PATENT

Attorney's Docket No.: P-9902 MG

DECLARATION FOR PATENT APPLICATION

As a below-named inventor, I hereby declare that:

My residence, post office address and citizenship are as stated below next to my name, I believe I am the original, first and sole inventor (if only one name is listed below) or an original, first and joint inventor (if plural names are listed below) of the subject matter which is claimed and for which a patent is sought on the invention entitled PAYMENT METHOD AND APPARATUS THEREFOR the specification of which

(check one) ☒ is attached hereto.

☐ was filed on _____ as _____

Application Serial No.:

and was amended on _____ (if applicable).

I hereby state that I have reviewed and understand the contents of the above-identified specification, including the claims, as amended by any amendment referred to above.

I acknowledge the duty to disclose information which is material to patentability as defined in Title 37, Code of Federal Regulations, §1.56(a).

I hereby claim foreign priority benefits under Title 35, United States Code, §119 of any foreign application(s) for patent or inventor's certificate listed below and have also identified below any foreign application for patent or inventor's certificate having a filing date before that of the application on which priority is claimed:

Prior Foreign Application(s):

Appln. No.	Country	Date Filed	Priority Claimed
98120922	Russia	November 25, 1998	<input checked="" type="checkbox"/> YES <input type="checkbox"/> NO
PCT/RU98/00264	PCT	July 29, 1999	<input checked="" type="checkbox"/> YES <input type="checkbox"/> NO
			<input type="checkbox"/> YES <input type="checkbox"/> NO

I hereby claim the benefit under Title 35, United States Code, §120 of any United States application(s) listed below and, insofar as the subject matter of each of the claims of this application is not disclosed in the prior United States application in the manner provided by the first paragraph of Title 35, United States Code §112, I acknowledge the duty to disclose information which is material to patentability as defined in Title 37, Code of Federal Regulations, §1.56(a) which became available between the filing date of the prior application and the national or PCT international filing date of this application:

Appln. Serial No.	Filing Date	Status: Patented, Pending, Abandoned
		<input type="checkbox"/> Patented <input type="checkbox"/> Pending <input type="checkbox"/> Abandoned
		<input type="checkbox"/> Patented <input type="checkbox"/> Pending <input type="checkbox"/> Abandoned
		<input type="checkbox"/> Patented <input type="checkbox"/> Pending <input type="checkbox"/> Abandoned

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POWER OF ATTORNEY

I hereby appoint the following attorney(s) and/or agent(s) to prosecute the application entitled **PAYMENT METHOD AND APPARATUS THEREFOR** and to transact all business in the Patent and Trademark Office connected therewith:

(4) **HENRY A. MARZULLO, JR., Reg. No. 20,910;**

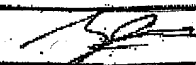

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I hereby declare that all statements made herein of my own knowledge are true and that all statements made on information and belief are believed to be true; and further that these statements were made with the knowledge that willful false statements and the like so made are punishable by fine or imprisonment, or both, under Section 1001 of Title 18 of the United States Code, and that such willful false statements may jeopardize the validity of the application or any patent issued thereon.

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Date Nov. 5, 1999	Signature 

Additional inventors are being named on separately numbered sheets attached hereto.

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RESIDENCE Address -- Street <u>ul. Babushkina, 29, korp. 2, kv. 45</u>	POST OFFICE Address -- Street (same as residence)
City (Zip) <u>193029, Saint-Petersburg</u>	City (Zip)
State or Country <u>Russia</u>	State or Country
Date <u>Nov. 5, 1999</u>	Signature <u>[Signature]</u>

☐ Additional inventors are being named on separately numbered sheets attached hereto.

914 723 4300 P.20/27

Docket No.: P-9902 MG

Page 1 of 2

For: Payment Method and Apparatus Therefor

[37 CFR 1.9(f) and 1.27(C)]

INDEPENDENT INVENTOR

As one of the five below named inventors,

Inventor	Address
Oleg Anatolovich ZOLOTOREV	Sosnovy bor, ul. Molodegnaya, 25, kv.20 188537, Leningradskaya obl. RUSSIA
Ivan Vladimirovich KUZNETSOV	ul. Saltikova-Chedrina, 48, kv. 56 191123, Saint-Petersburg RUSSIA
Andrey Gennadievich MOSEONKIN	ul. Shpalernaya, 44-A, kv. 5 191123, Saint-Petersburg RUSSIA
Alexandr Leonidovich SMIRNOV	ul. Dostoevskogo, 36, kv. 8 191126, Saint-Petersburg RUSSIA
Ildar Magafurovich KHAMITOV	ul. Babushkina, 29, korp. 2, kv. 45 193029, Saint-Petersburg RUSSIA

I hereby declare that I qualify as an independent inventor as defined in 37 CFR 1.9(c) for purposes of paying reduced fees under section 41(a) of Title 35, United States Code, to the Patent and Trademark Office with regard to the invention entitled **Payment Method and Apparatus Therefor** described in

- ☒ the specification filed herewith
☐ application serial no. _____, filed on _____
☐ Patent No. _____, issued on _____

I have not assigned, granted, conveyed or licensed and am under no obligation under contract or law to assign, grant, convey or license, any rights in the invention to any person who could not be classified as an independent inventor under 37 CFR 1.9(c) if that person had made the invention, or to any concern which would not qualify as a small business concern under 37 CFR 19(d) or a nonprofit organization under 37 CFR 1.9(e).

Each person, concern or organization to which I have assigned, granted, conveyed, or licensed or am under an obligation under contract or law to assign, grant, convey, or license any rights in the invention is listed below:

- ☐ no such person, concern, or organization
☐ persons, concerns, or organizations listed below.*

Q. Was the witness ever in the company of the defendant, or of any of the persons named in the indictment, at the time of the commission of the crime?

NOV-03-1999 11:17

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914 723 4300 P.21/27

New National Phase PCT Application in the U.S. for "Payment Method and Apparatus Therefor," Page 2 of 2

Note: Separate verified statements are required from each named person, concern or organization having rights to the invention averring to their status as small entities. (37 CFR 1.27)


FULL NAME: ADDRESS:	<input type="checkbox"/> INDIVIDUAL <input type="checkbox"/> SMALL BUSINESS CONCERN <input type="checkbox"/> NONPROFIT ORGANIZATION
FULL NAME: ADDRESS:	<input type="checkbox"/> INDIVIDUAL <input type="checkbox"/> SMALL BUSINESS CONCERN <input type="checkbox"/> NONPROFIT ORGANIZATION
FULL NAME: ADDRESS:	<input type="checkbox"/> INDIVIDUAL <input type="checkbox"/> SMALL BUSINESS CONCERN <input type="checkbox"/> NONPROFIT ORGANIZATION

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I hereby declare that all statements made herein of my own knowledge are true and that all statements made on information and belief are believed to be true; and further that these statements were made with the knowledge that willful false statements and the like so made are punishable by fine or imprisonment, or both, under Section 1001 of Title 18 of the United States Code, and that such willful false statements may jeopardize the validity of the application, any patent issuing thereon, or any patent to which this verified declaration is directed.

NAME OF INVENTOR

Ivan Vladimirovich KUZNETSOV

SIGNATURE OF INVENTOR**DATE**

Nov 5, 1999

LACKENBACH SIEGEL MARZULLO ARONSON & GREENSPAN, P.C. SE-B-1

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914 723 4300 P.23/27

New National Phase PCT Application in the U.S. for "Payment Method and Apparatus Therefor," Page 2 of 2

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NAME OF INVENTOR

Andrey Gennadievich MOSHONKIN

SIGNATURE OF INVENTOR



DATE

Nov 5, 1999

LACKENBACH SIEGEL MARZULLO ARONSON & GREENSPAN, P.C. SE-B-1

663027-9854460

NOV-03-1999 11:15 LACKENBACH SIEGEL
Applicant or Patentee: **Zolotarev et al.**

914 723 4300 P.18/27
Docket No.: **P-9902 MG**

Serial or Patent No.: (to be assigned)

Page 1 of 2

Filed or Issued:

For: **Payment Method and Apparatus Therefor**

DECLARATION CLAIMING SMALL ENTITY STATUS

(37 CFR 1.9(f) and 1.27(c))

INDEPENDENT INVENTOR

As one of the five below named inventors,

Inventor	Address
Oleg Anatolovich ZOLOTAREV	Sosnovy bor, ul. Molodegnaya, 25, kv.20 188537, Leningradskaya obl. RUSSIA
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914 723 4320 P.19/27

New National Phase PCT Application in the U.S. for "Payment Method and Apparatus Therefor," Page 2 of 2

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NAME OF INVENTOR

Oleg Anatolovich ZOLOTOREV

SIGNATURE OF INVENTOR**DATE**

11.5.99

LACKENBACH SIEGEL MARZULLO ARONSON & GREENSPAN, P.C. SE-B-1

663021 33E54460

NOV-03-1999 11:19

LACKENBACH SIEGEL

914 723 4300 P.24/27

Applicant or Patentee: **Zolotarev et al.**Docket No.: **P-9902 MG**

Serial or Patent No.: (to be assigned)

Page 1 of 2

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For: **Payment Method and Apparatus Therefor****DECLARATION CLAIMING SMALL ENTITY STATUS**

[37 CFR 1.9(c) and 1.27(c)]

INDEPENDENT INVENTOR

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669021-88E54460

NOV-03-1999 11:20
Applicant or Patentee:LACKENBACH SIEGEL
ZOLOTOREV et al.914 723 4300 P.26/27
Docket No.: P-8902 MG

Serial or Patent No.: (to be assigned)

Page 1 of 2

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For: Payment Method and Apparatus Therefor

DECLARATION CLAIMING SMALL ENTITY STATUS

[37 CFR 1.9(c) and 1.27(C)]

INDEPENDENT INVENTOR

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669021-88E54450

NOV-23-1999 11:21 LACKENBACH SIEGEL 914 723 4300 P.27/27
new national phase PCT Application in the U.S. for "Payment Method and Apparatus
Therefor," Page 2 of 2

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NAME OF INVENTOR

Ildar Magafurovich KHAMITOV

SIGNATURE OF INVENTOR

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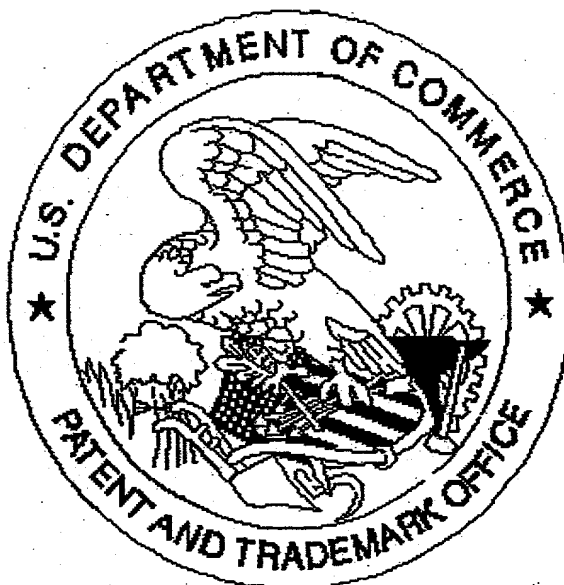
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